

# Multi-agent techniques for resource allocation and planning

## Application to Earth observation by satellite constellations

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ENS Rennes – 07/12/2023

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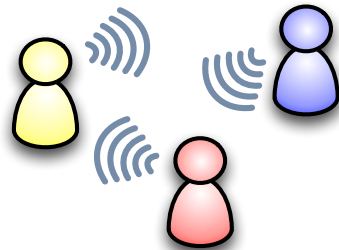
# Introduction

## Multi-Agent Systems and Distributed Artificial Intelligence

- **Agent:** An entity that behaves autonomously in the pursuit of goals
- **Multi-agent system:** A system of multiple interacting agents

### An agent is...

- **Autonomous:** Is of full control of itself
- **Interactive:** May communicate with other agents
- **Reactive:** Responds to changes in the environment or requests by other agents
- **Proactive:** Takes initiatives to achieve its goals



# Introduction

## Sample multi-agent systems



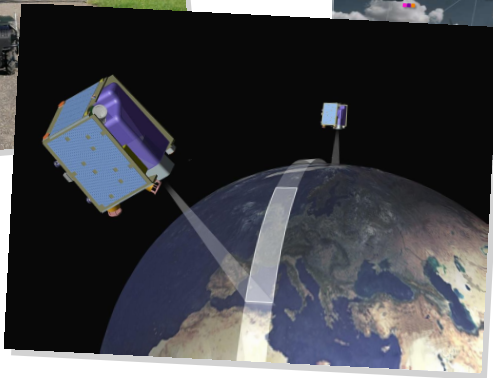
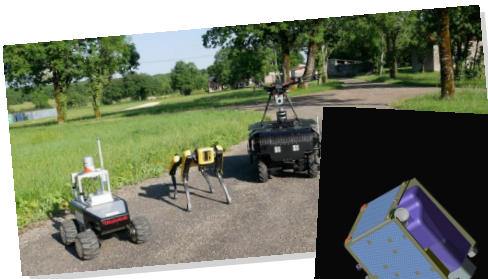
# Introduction

## Sample multi-agent systems



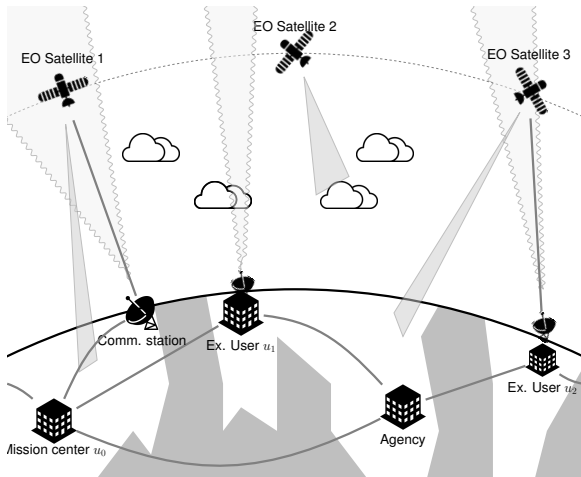
# Introduction

## Sample multi-agent systems



# Introduction

## Multi-Agent Decisions for Earth Observation



### • Constellation Design

- How to compose the constellation?
- How to dimension the constellation?
- Where to position assets?

### • Offline Operations

- How to allocate resources?
- How to share resources?
- How to schedule in a multi-party/multi-mission context?

### • Online Operations

- How to adapt activities facing unpredictable events?
- Which coordination protocols to use?

# Outline

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- 1 Introduction
- 2 Challenges in Earth Observation Constellation Operations
- 3 Focus #1: Sharing Space Assets
- 4 Focus #2: Coordinating Asset Usage
- 5 Conclusion

# Outline

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- 1 Introduction
- 2 Challenges in Earth Observation Constellation Operations**
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# Challenges in Earth Observation Constellation Operations

[PICARD et al., 2021]

- Recent years have shown a large increase in the development of satellite constellations
- Increasing the size allows to capture any point on Earth at higher frequency, e.g. the Planet Dove constellation
- But, operating numerous Earth observation satellites (**EOS**) requires to *cooperate*, *collectively* solve and schedule, *self*-adapt and *interact*

Many AAMAS-related and Open Research and Technology Questions

# Categories of Challenges

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Constellation Design

Offline Operations

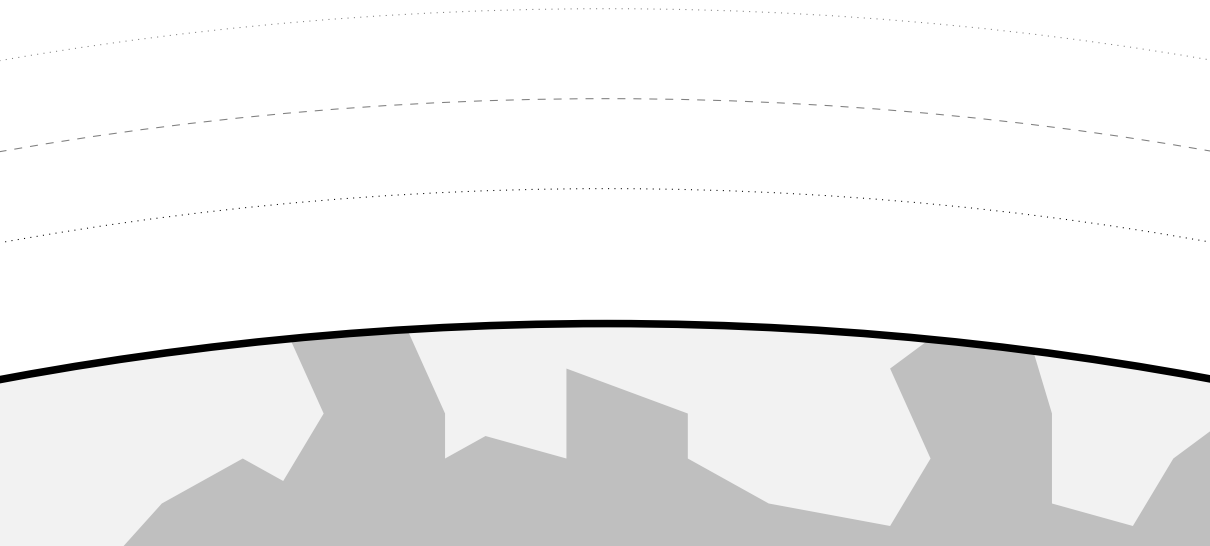
Online Operations

# How to Design an EOS Constellation?

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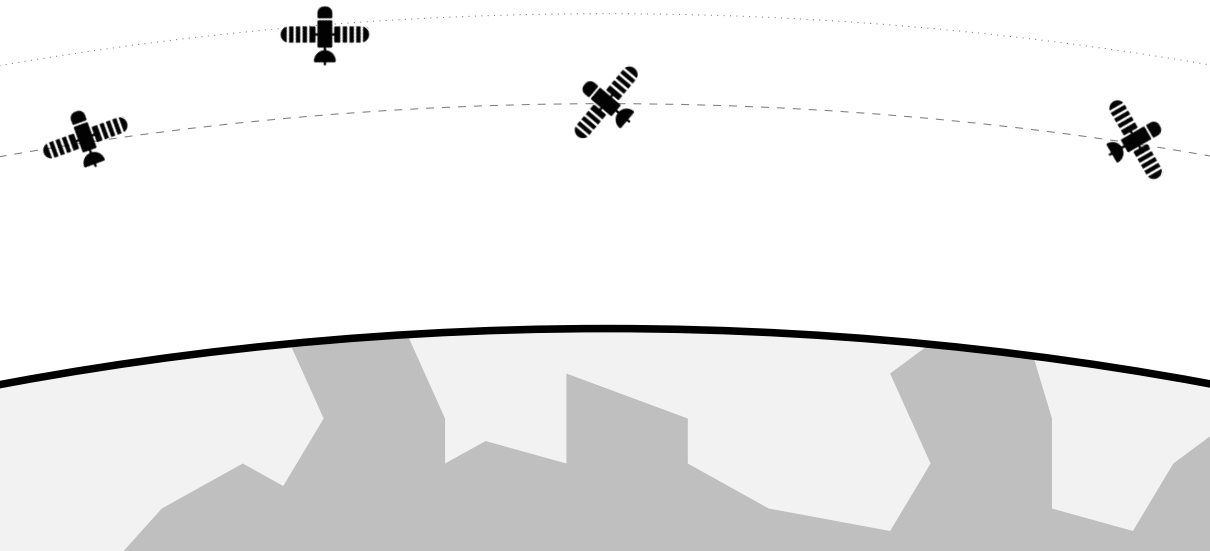
# How to Design an EOS Constellation?

## Orbits



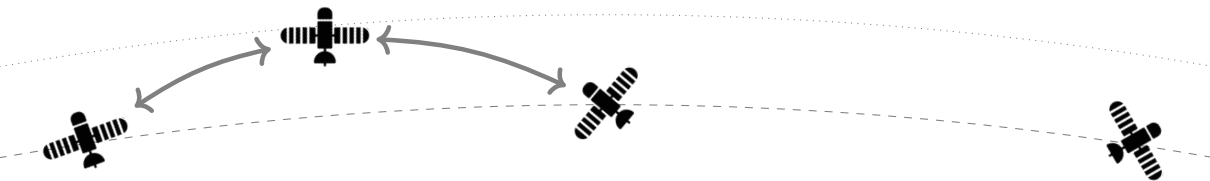
# How to Design an EOS Constellation?

## Constellation composition



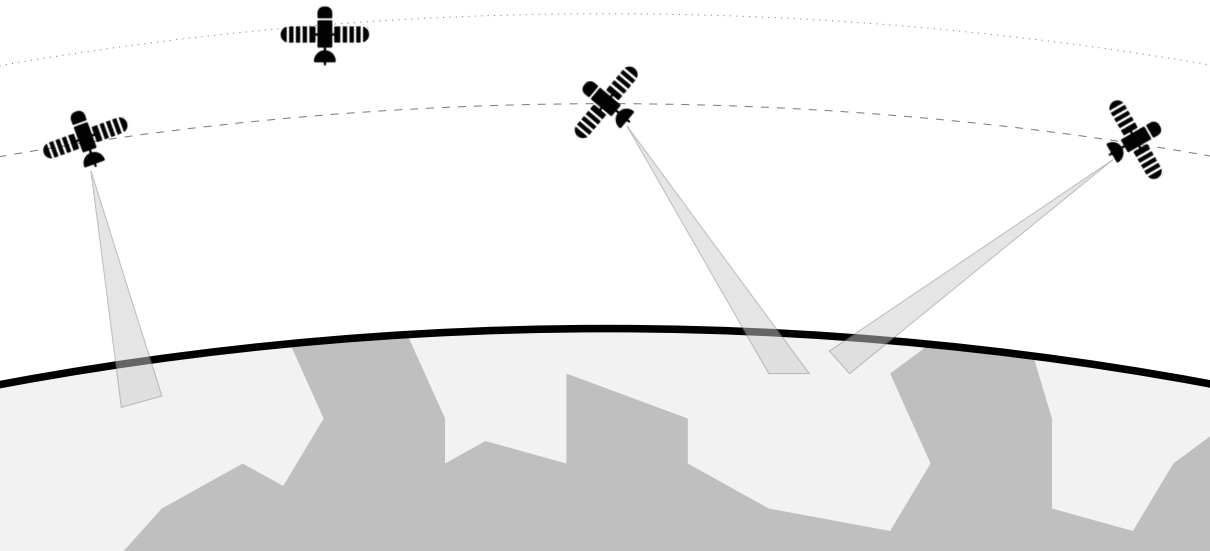
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## Constellation composition



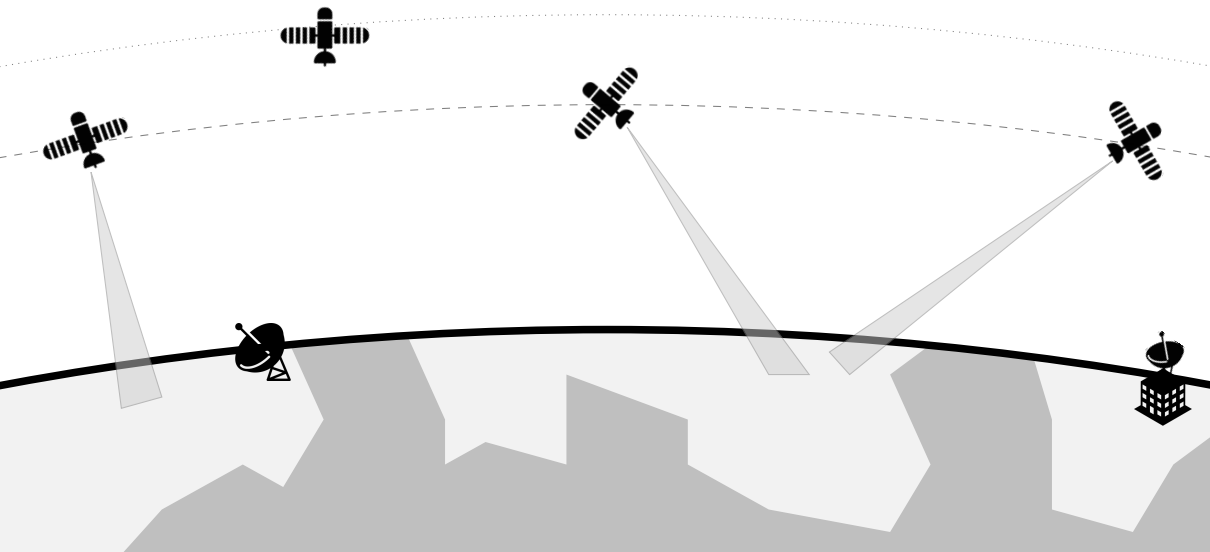
# How to Design an EOS Constellation?

Points of interest



# How to Design an EOS Constellation?

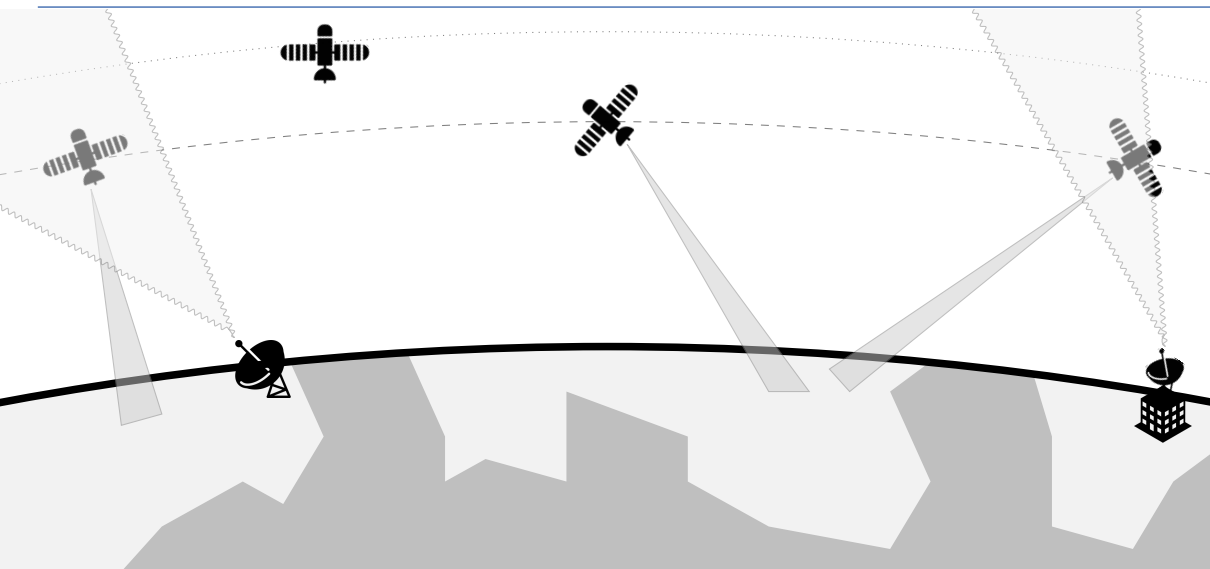
On-ground communication stations





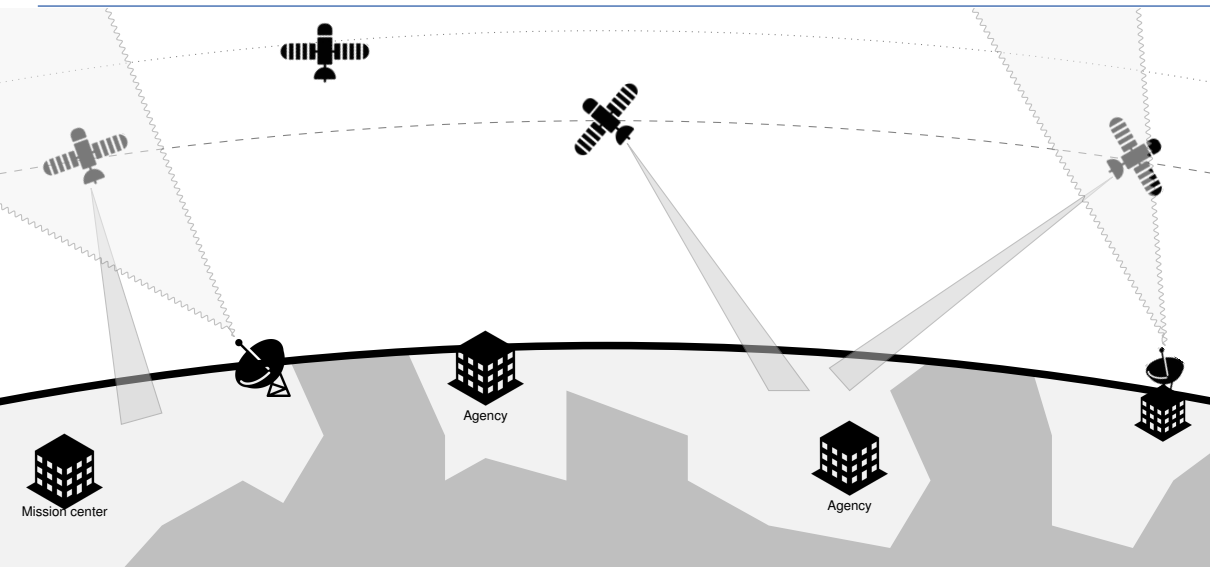
# How to Design an EOS Constellation?

## Visibility windows



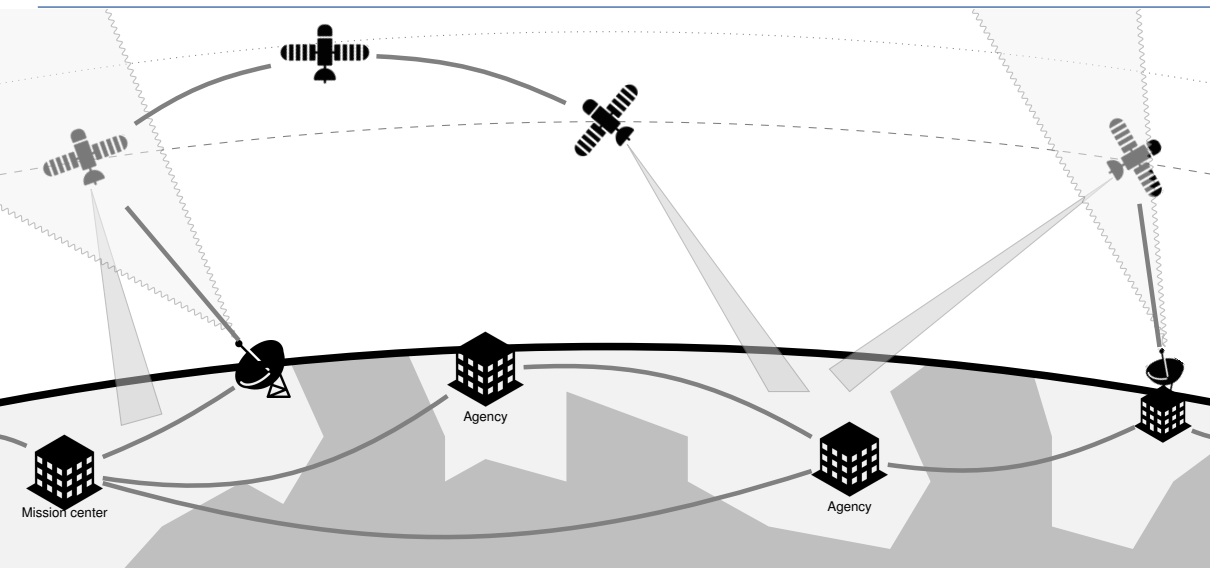
# How to Design an EOS Constellation?

## Other actors and stakeholders



# How to Design an EOS Constellation?

## System organization

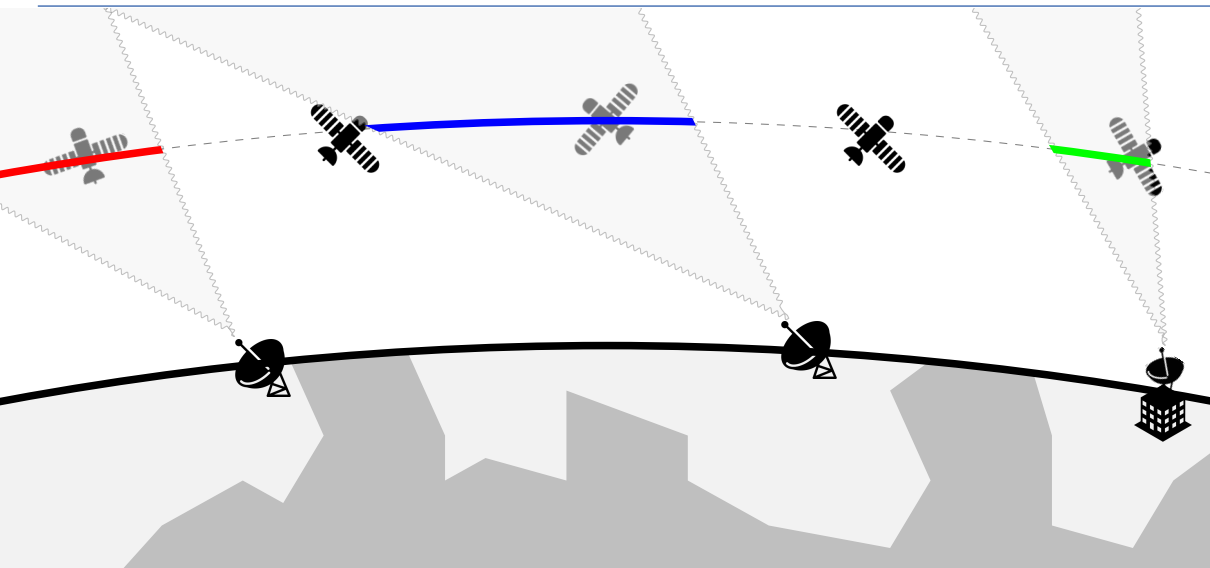


# How to Allocate Resources?

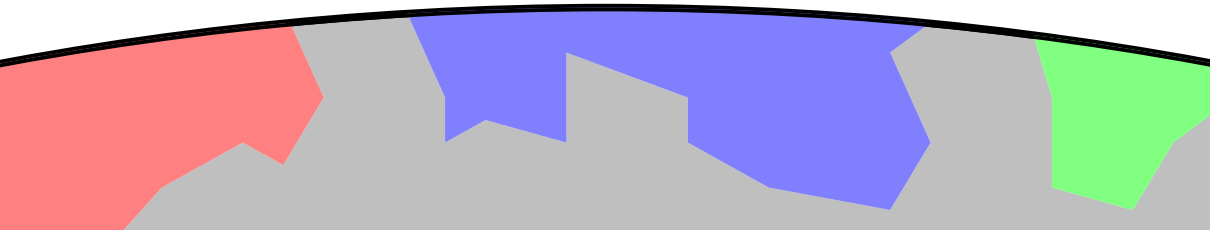
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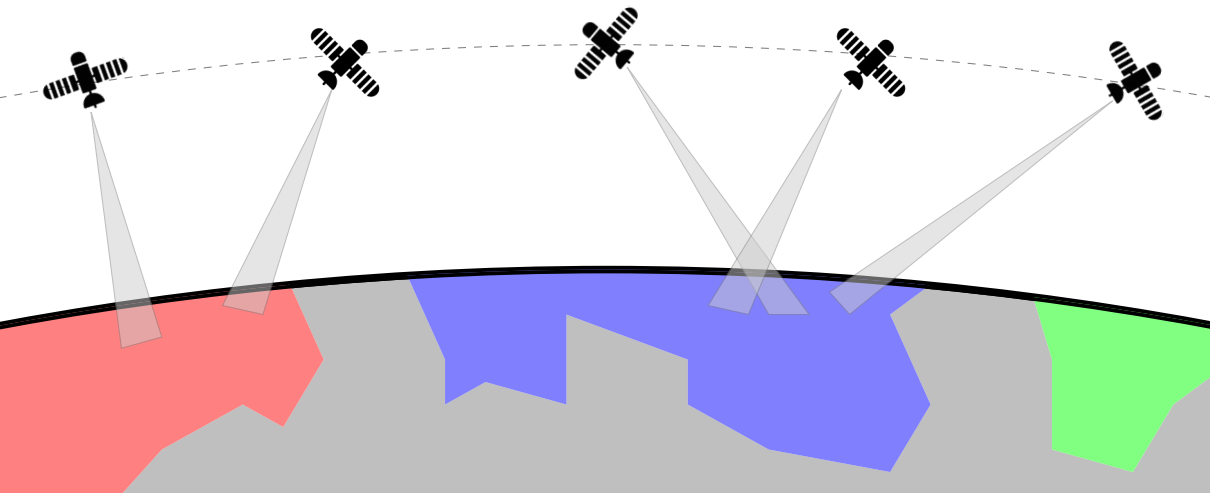
# How to Allocate Resources?



# How to Share Resources?



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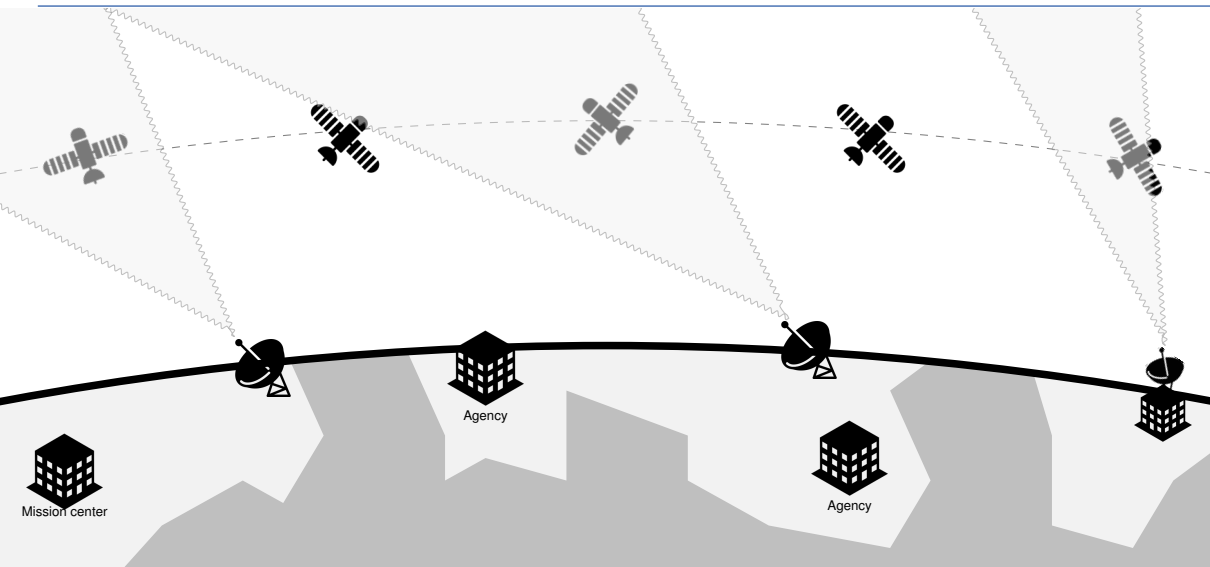


# How to Schedule in a Multi-satellite and Multi-user Setting?

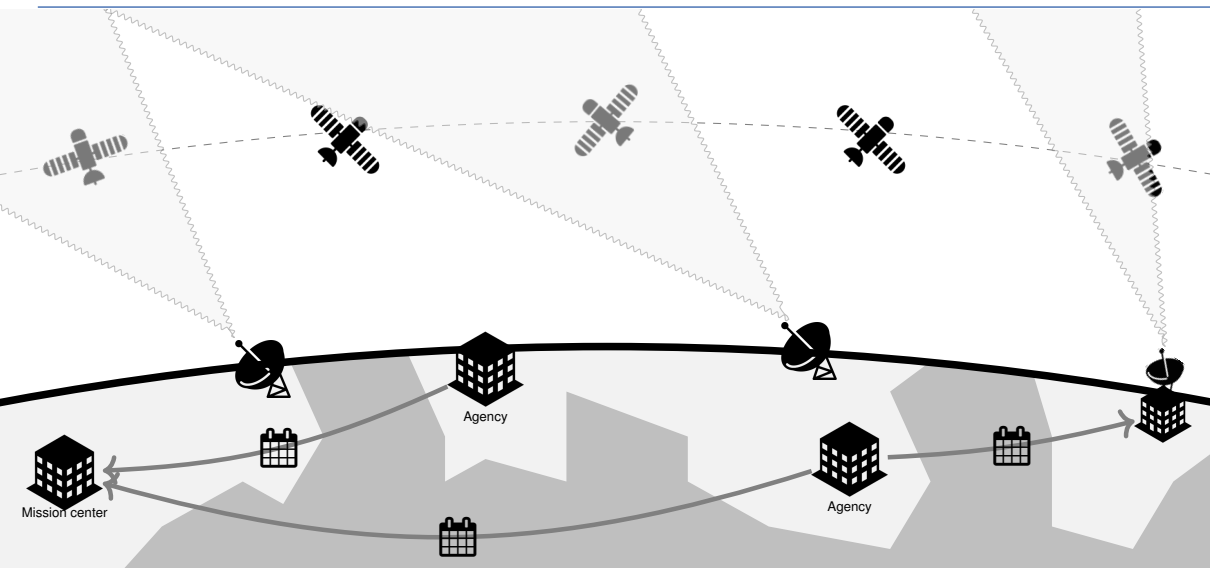




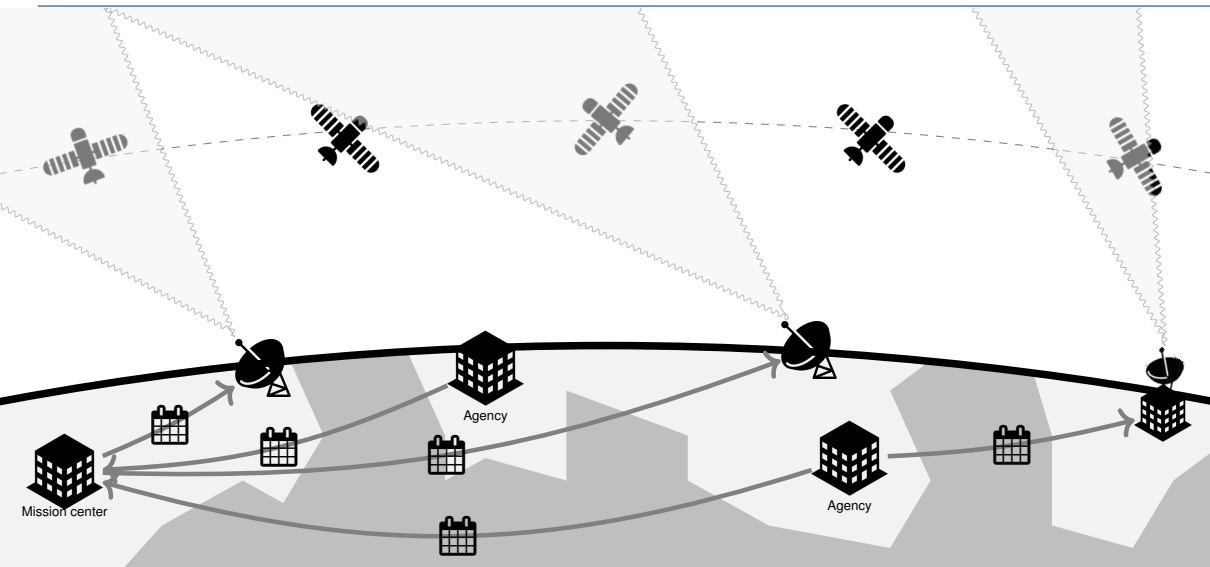
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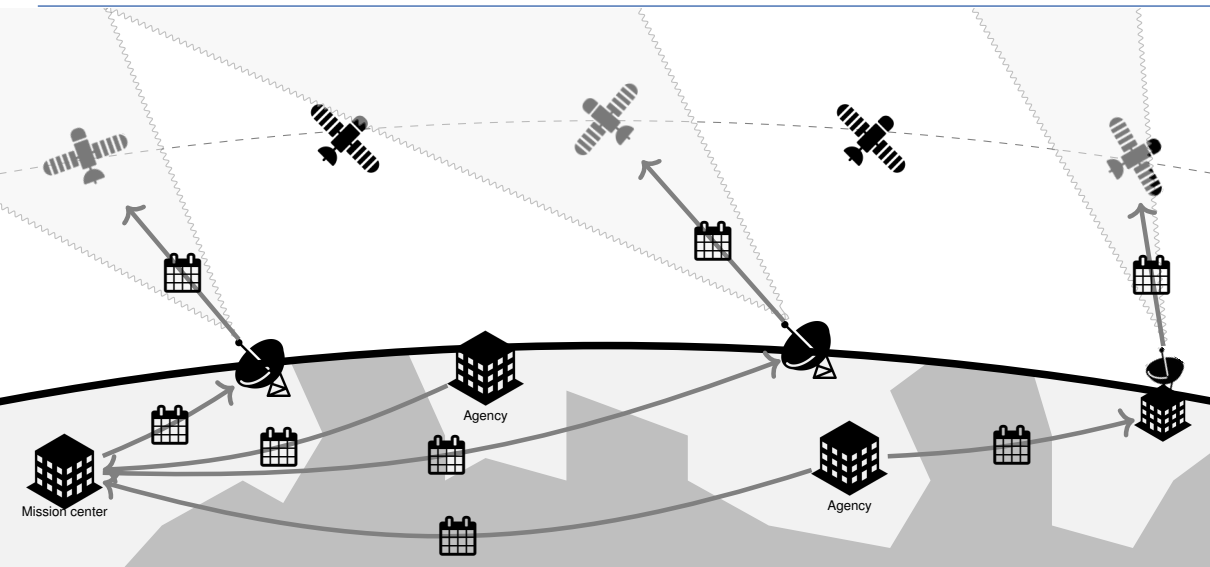
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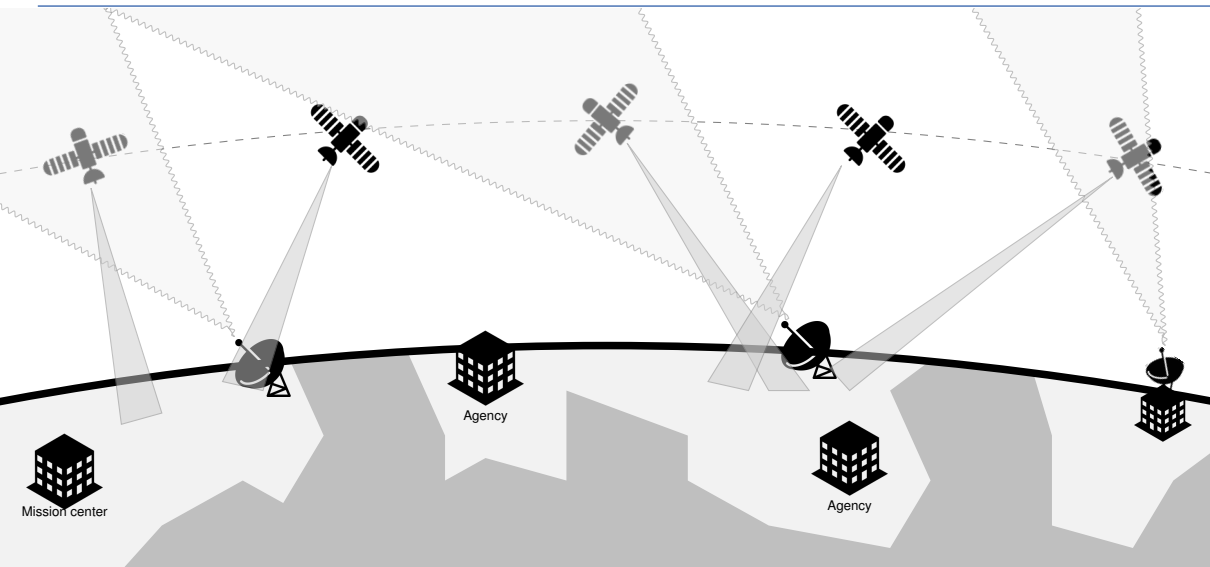
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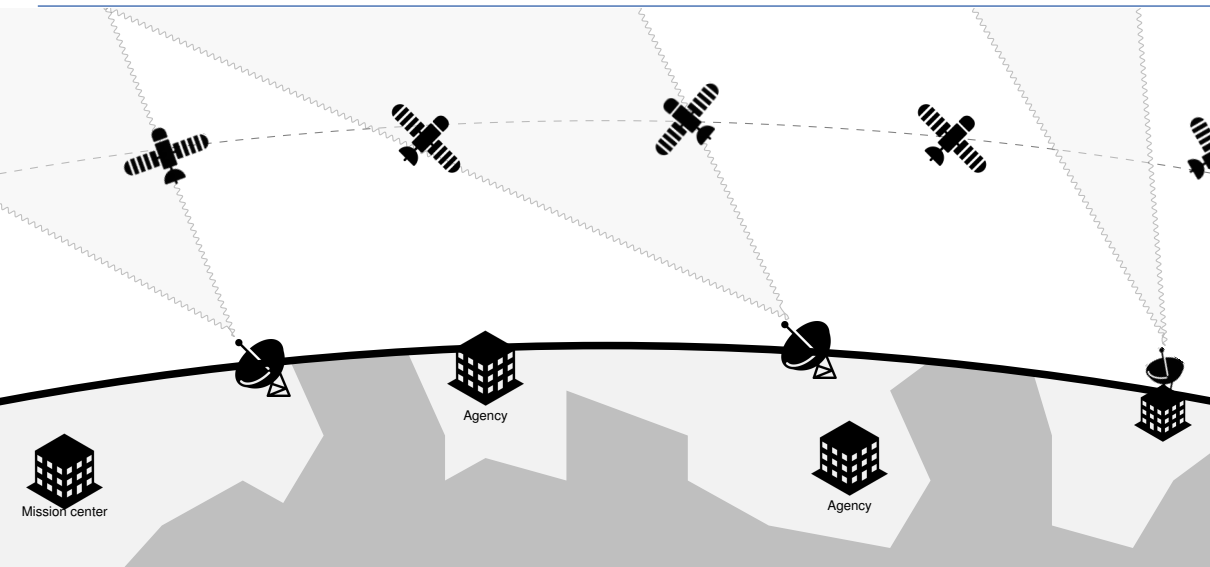
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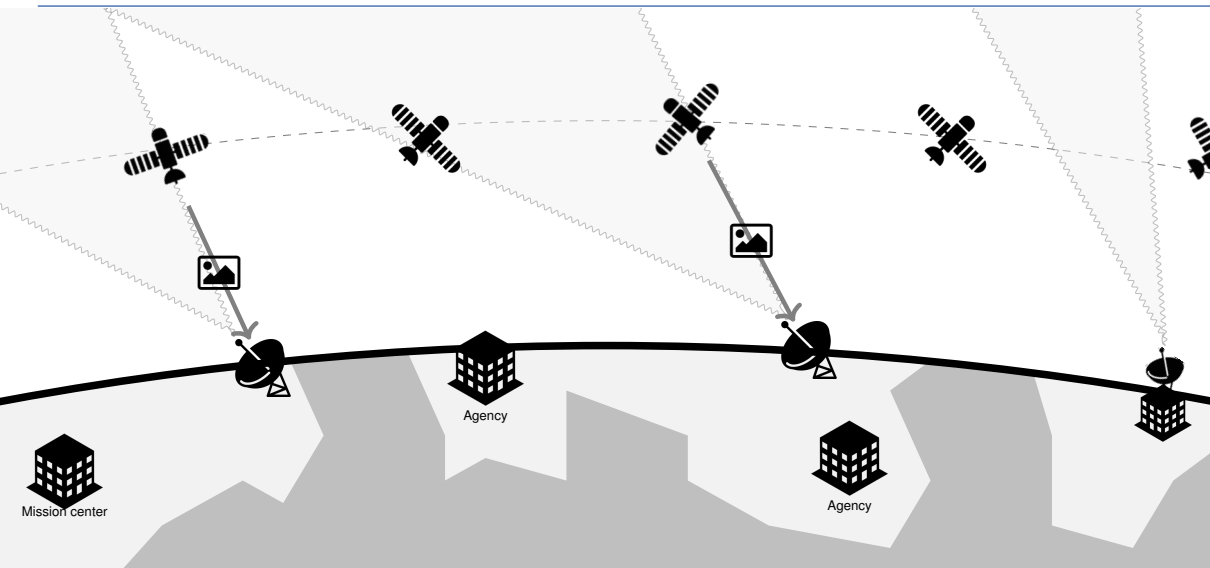
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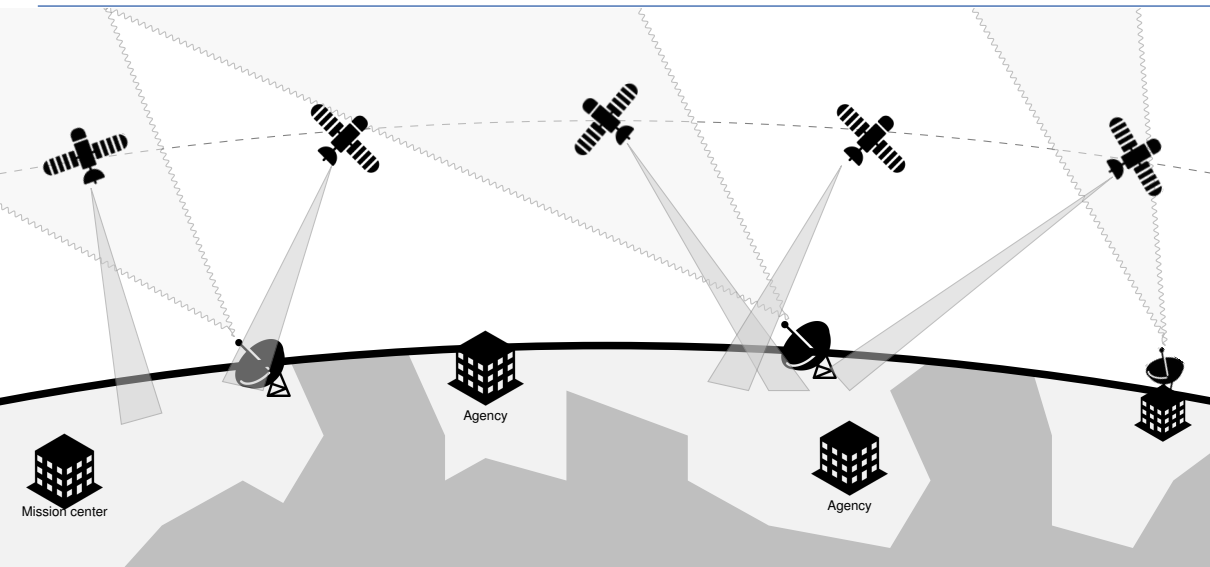
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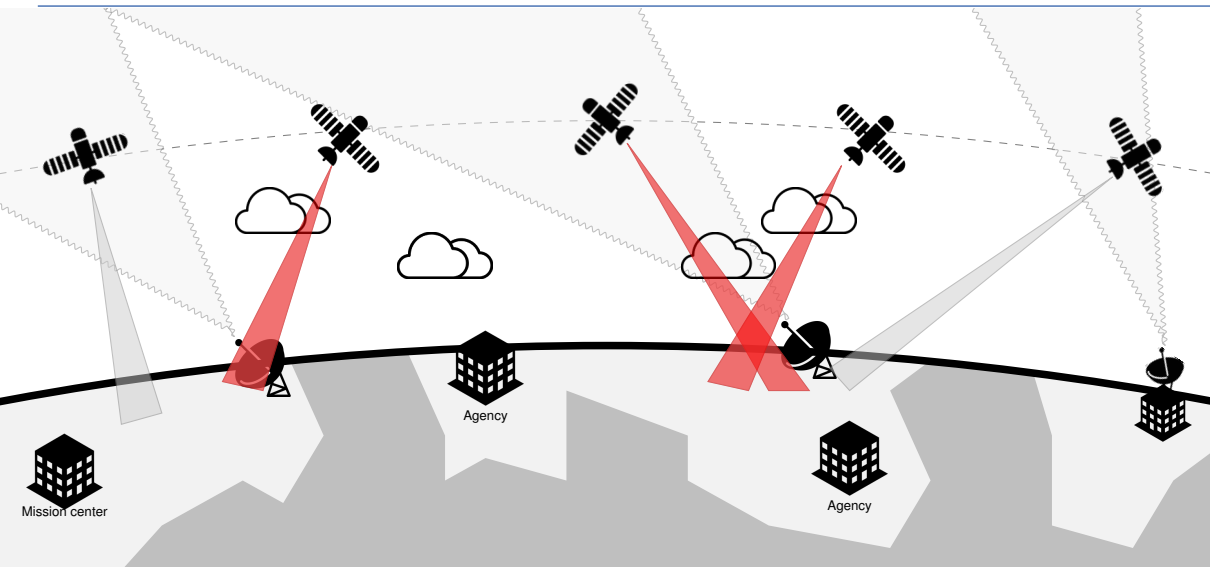


# How to Adapt Activities when Facing Unpredictable Events?

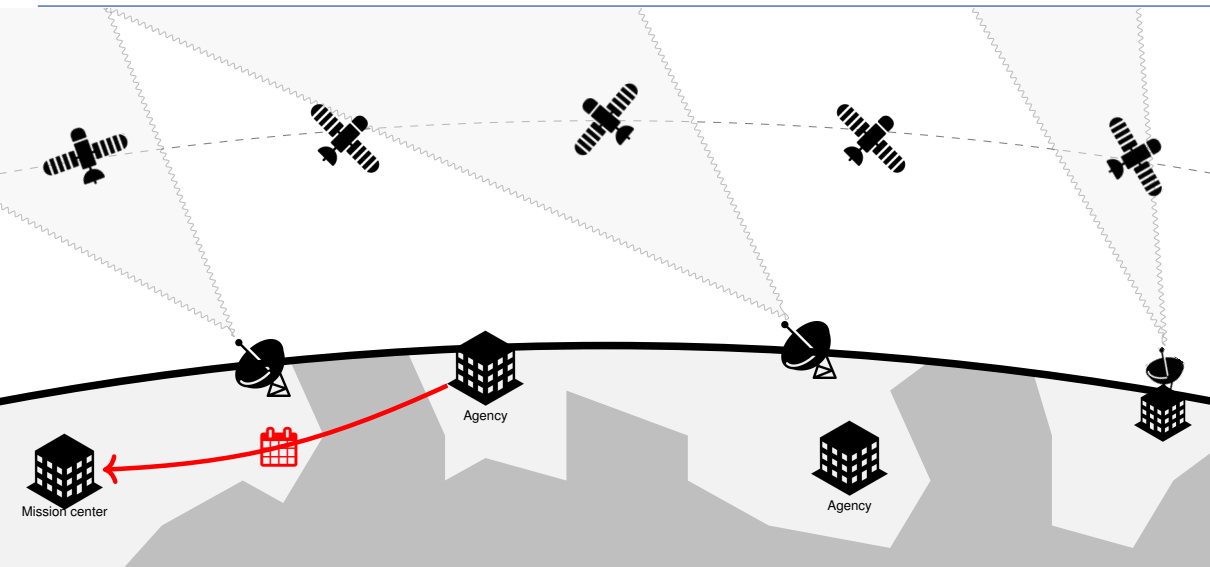




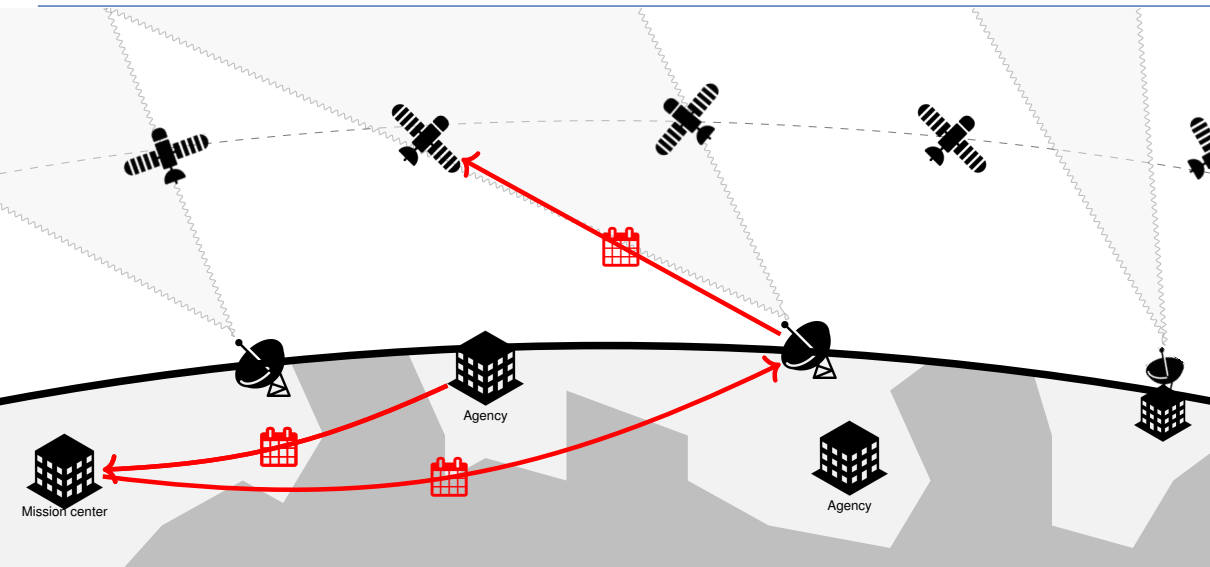
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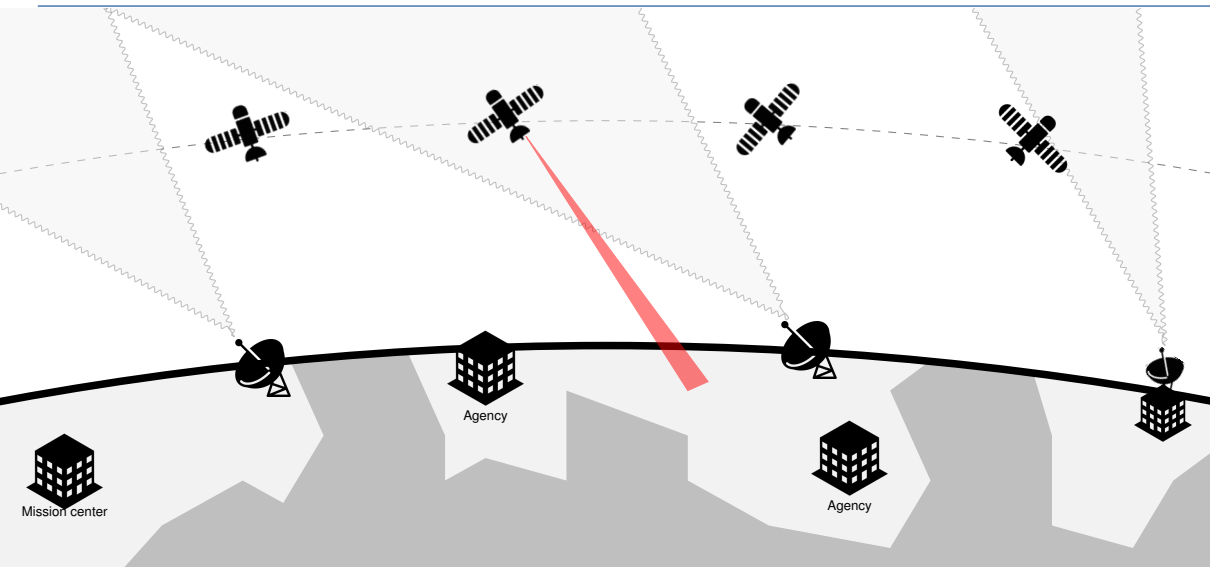
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# Outline

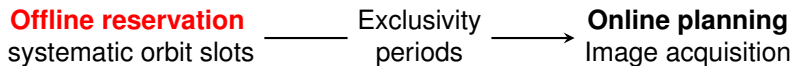
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# Sharing Space Assets

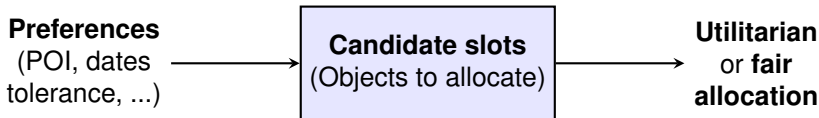
Example: Earth observation satellite constellations

- **Problem** : exploitation of the same constellation/mission by several stakeholders



- **Current allocation scheme**: first come, first served

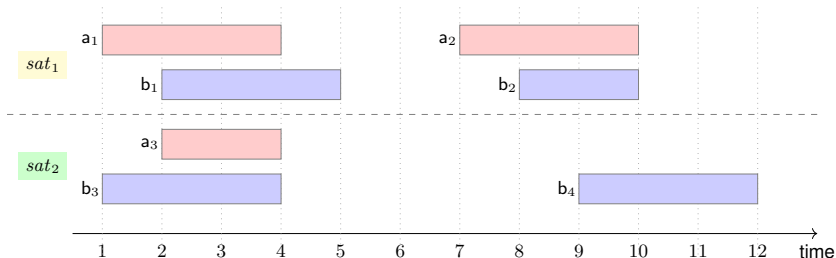
- **Objective**



# Orbit Slot Allocation

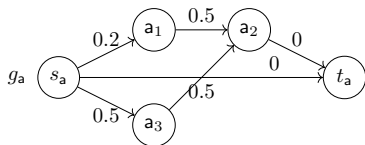
## An example

- 2 agents ( $a$  in red,  $b$  in blue) requesting acquisitions:
  - of points of interest (POI) around the same region
  - around 2 time points (3 and 9) every day
- 1 satellite giving access to 2 orbit slots for each time point ( $a_1, \dots, a_3, b_1, \dots, b_4$ )



# Orbit Slot Allocation Problem

## Graph representation

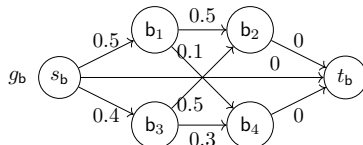


Paths for graph  $g_a$ :

$$\pi_{a,0} = [s_a, t_a]$$

$$\pi_{a,1} = [s_a, a_1, a_2, t_a]$$

$$\pi_{a,2} = [s_a, a_3, a_2, t_a]$$



Paths for graph  $g_b$ :

$$\pi_{b,0} = [s_b, t_b]$$

$$\pi_{b,1} = [s_b, b_1, b_2, t_b]$$

$$\pi_{b,2} = [s_b, b_1, b_4, t_b]$$

$$\pi_{b,3} = [s_b, b_3, b_2, t_b]$$

$$\pi_{b,4} = [s_b, b_3, b_4, t_b]$$

Forbidden combinations:

$$(\pi_{a,1}, \pi_{b,1})$$

$$(\pi_{a,1}, \pi_{b,3})$$

$$(\pi_{a,2}, \pi_{b,1})$$

$$(\pi_{a,2}, \pi_{b,3})$$

$$(\pi_{a,2}, \pi_{b,4})$$



# Problem Model

## The non-compact case

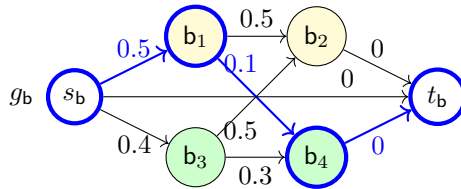
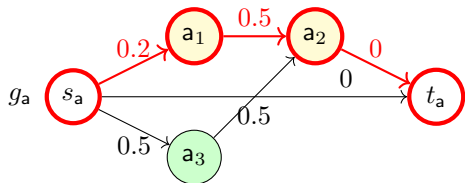
### Definition

A *Directed Path Allocation Problem* (DPAP) is a tuple  $\langle \mathcal{A}, \mathcal{G}, \mu, \phi \rangle$ , where

- $\mathcal{A} = \{1, \dots, n\}$  is a set of agents
- $\mathcal{G} = \{g_1, \dots, g_m\}$  is a set of single-source single-sink edge-weighted DAGs
- $\mu : \mathcal{G} \rightarrow \mathcal{A}$  maps each graph  $g$  in  $\mathcal{G}$  to its owner  $a$  in  $\mathcal{A}$ ; we also denote by  $\mathcal{G}_a = \mu^{-1}(a)$  the set of graphs owned by agent  $a$
- $\phi : \Pi_{g_1} \times \dots \times \Pi_{g_m} \rightarrow \{0, 1\}$  is a *path compatibility function* that indicates whether a combination of paths  $(p_1, \dots, p_m)$  (one path per graph) is feasible (value 1) or not (value 0)

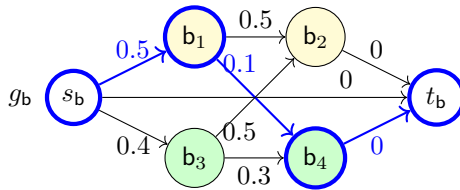
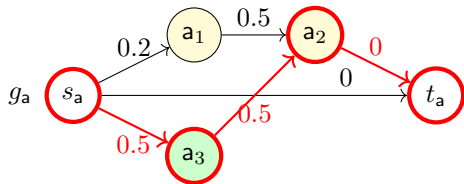
# DPAP Solutions

Selecting non conflicting path in each graph



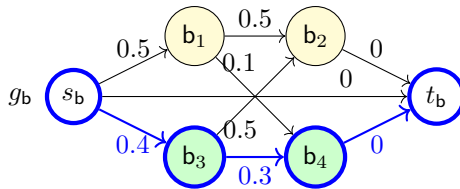
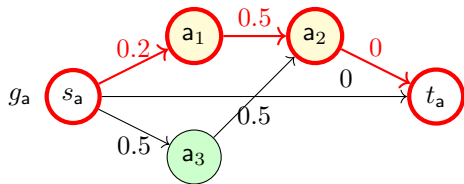
# DPAP Solutions

Selecting non conflicting path in each graph, **maximizing global utility**



# DPAP Solutions

Selecting non conflicting path in each graph, **maximizing fairness**

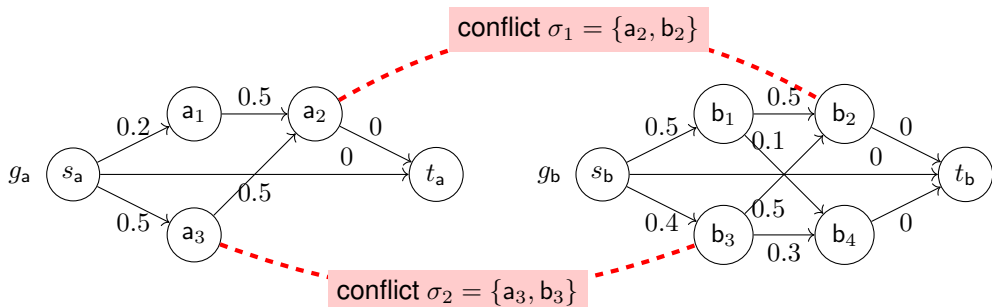


# DPAP Conflict Formulations

More compact ways to represent conflicts

## V-DPAP: Vertex-constrained Directed Path Allocation Problems

- $\phi$  is defined by a set of conflicts  $\mathcal{C}$  between vertices of the graph
- each conflict  $\sigma \in \mathcal{C}$  is a non-empty set of vertices  $V_\sigma$  that cannot be all selected by an allocation

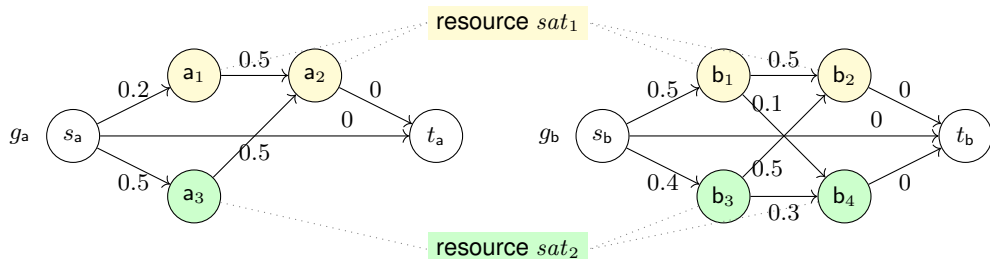


# DPAP Conflict Formulations (cont.)

More compact ways to represent conflicts

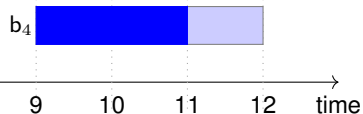
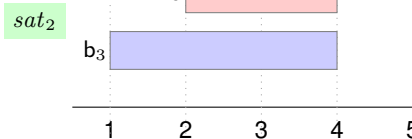
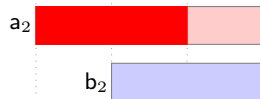
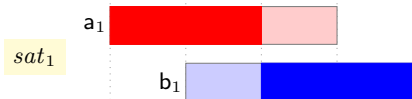
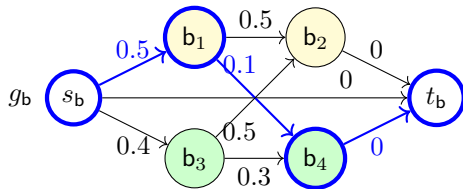
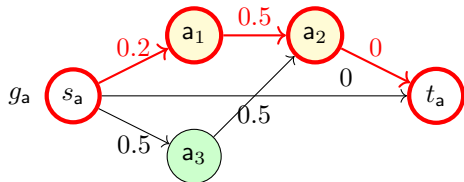
## R-DPAP: Resource-constrained Directed Path Allocation Problems

- $\phi$  considers a set of disjunctive resources  $\mathcal{R} = \{r_1, \dots, r_p\}$
- each vertex in the graph has start date, an end date, a duration, and a required resource
- there is a conflict if at least two time windows overlap on the same resource when scheduling without any interruption (non preemptive consumption)



# DPAP Conflict Formulations (cont.)

More compact ways to represent conflicts



1 2 3 4 5 6 7 8 9 10 11 12 time

# Properties

---

- V-DPAP is *NP-complete* (via reduction of 3-SAT)
- R-DPAP is *NP-complete* (via reduction of 1-machine scheduling problem)
- There exists an equivalent V-DPAP to any R-DPAP
  - by generating a set of item selection conflicts that is equivalent to the set of selections forbidden by the scheduling problem



# Properties

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- V-DPAP is *NP-complete* (via reduction of 3-SAT)
  - R-DPAP is *NP-complete* (via reduction of 1-machine scheduling problem)
  - There exists an equivalent V-DPAP to any R-DPAP
    - by generating a set of item selection conflicts that is equivalent to the set of selections forbidden by the scheduling problem
- We focus on the definition of algorithms for solving V-DPAP (because limited number of requests)

# How to solve V-DPAP?

Sorry, no detail here... see [MAQROT et al., 2022; ROUSSEL et al., 2023b]

1	Optimal utilitarian allocation (util)	MILP-based
2	Optimal leximin allocation (lex)	MILP-based iterated w/ revision
3	Approximate leximin allocation (a-lex)	MILP-based iterated wo/ revision
4	Greedy allocation (greedy)	adhoc
5	<i>round-robin</i> path allocation (p-rr)	adhoc
6	<i>round-robin</i> node allocation (n-rr)	adhoc

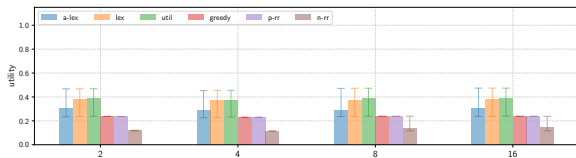
# Experimental Evaluation

Generation Parameters		Values
Constellation	Altitude	500 km
	Number of orbital planes $n_p$	2, 4, 8, 16
	Number of satellites/plane	2
	Inclination	40°
Scheduling horizon	Start	01-01-2020
	Duration	180 days
Problems	Number of users	4
	Type	V-DPAP, R-DPAP
Requests	Number of requests/user	2
	Requested Observation Times	3 RTs/request
	Maximum random time shift $\delta_r$	1 hour
	Tolerance $\Delta$	1 hour
	Minimum slot duration $minD$	120 seconds
Algorithms	Satisfaction mode	full, partial
	Type	util, lex, a-lex, greedy, p-rr, n-rr
	CPLEX Time Limit	120 seconds

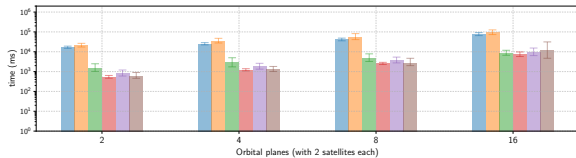
## Experimental Evaluation (cont.)

Problem	Properties	$n_p$			
		2	4	8	16
V-DPAP	Conflicts	37715.34	74009.12	146657.94	291831.52
	Conflict size	2.0	2.0	2.0	2.0
	Slots per RT	1.94	3.81	7.54	15.01
	Slot duration (s)	618.10	616.44	616.91	616.66
R-DPAP	Conflicts	1715.38	3527.42	6981.19	13929.55
	Conflict size	3.28	3.17	3.21	3.19
	Slots per RT	1.94	3.81	7.54	15.01
	Slot duration (s)	618.10	616.44	616.91	616.66

# Results for full request satisfaction mode

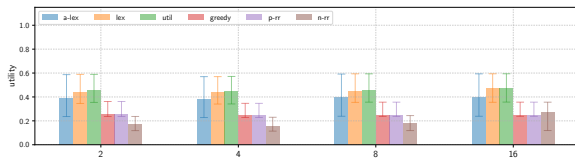


(a) Normalized utility

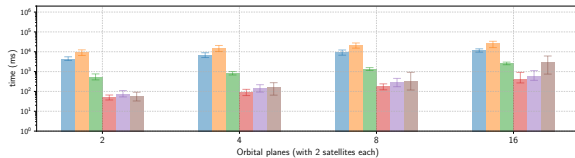


(b) Computation time

Figure: V-DPAP



(a) Normalized utility



(b) Computation time

Figure: R-DPAP

# Results for full request satisfaction mode (cont.)

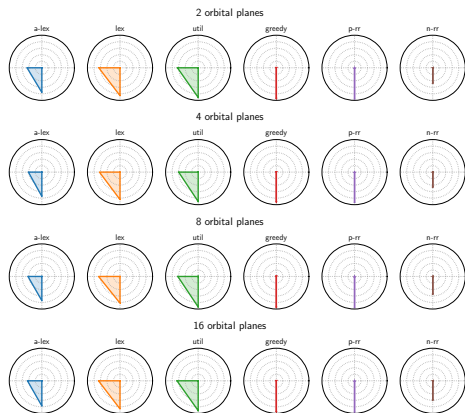


Figure: V-DPAP

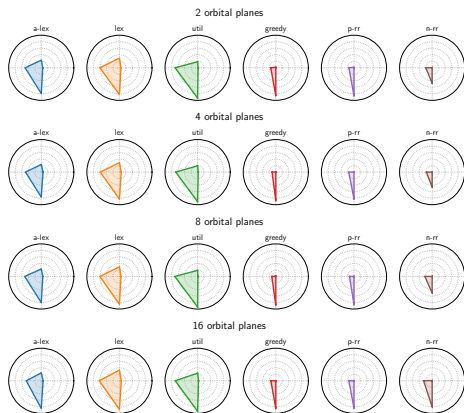
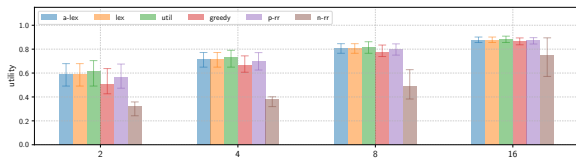
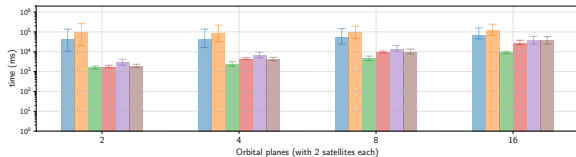


Figure: R-DPAP

# Results for flexible request satisfaction mode

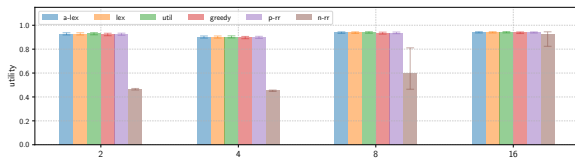


(a) Normalized utility

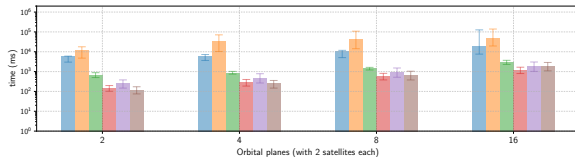


(b) Computation time

Figure: V-DPAP



(a) Normalized utility



(b) Computation time

Figure: R-DPAP

# Results for flexible request satisfaction mode (cont.)

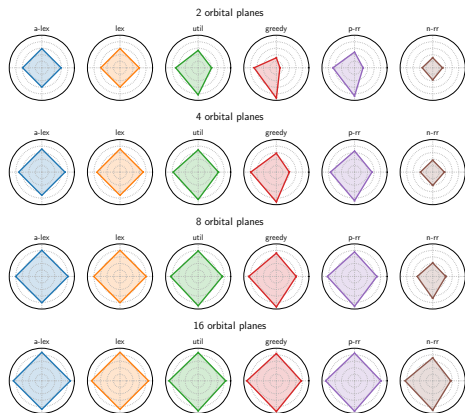


Figure: V-DPAP

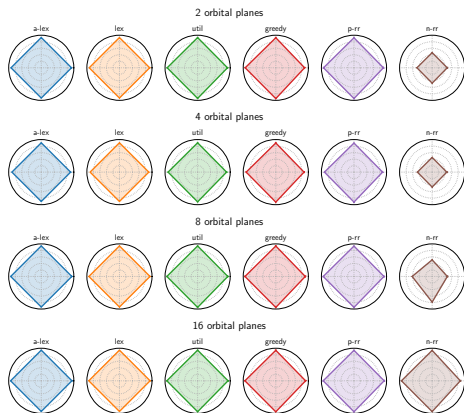
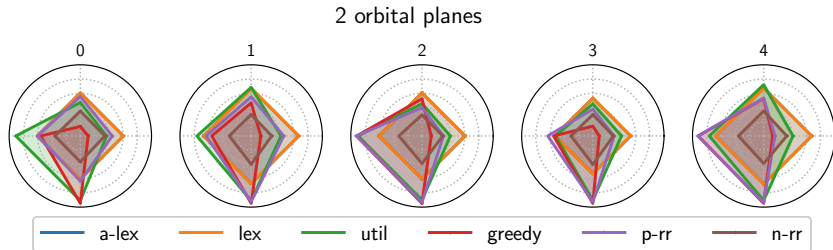


Figure: R-DPAP



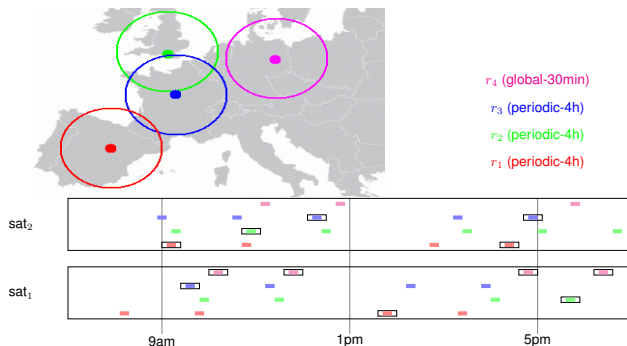
## Results for flexible request satisfaction mode (cont.)



**Figure:** Utility profiles (in leximin order) for the first 5 instances for a constellation with 2 orbital planes (4 satellites) and each algorithm (south: best utility over all agents; west: second best utility; north: third best utility; east: worst utility), for flexible requests encoded as V-DPAP.

# Where to find detailed info?

- Path allocation [MAQROT et al., 2022]
- DPAP and related methods [ROUSSEL et al., 2023b]
- More complex requests and CP-based methods [MAQROT et al., 2022]
- Some data [ROUSSEL et al., 2023a]



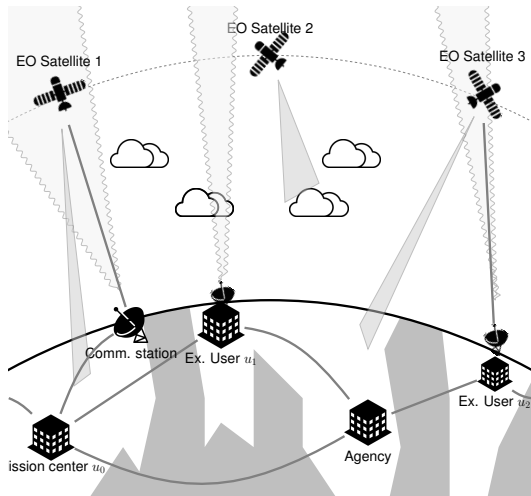
# Outline

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- 1 Introduction
- 2 Challenges in Earth Observation Constellation Operations
- 3 Focus #1: Sharing Space Assets
- 4 Focus #2: Coordinating Asset Usage**
- 5 Conclusion

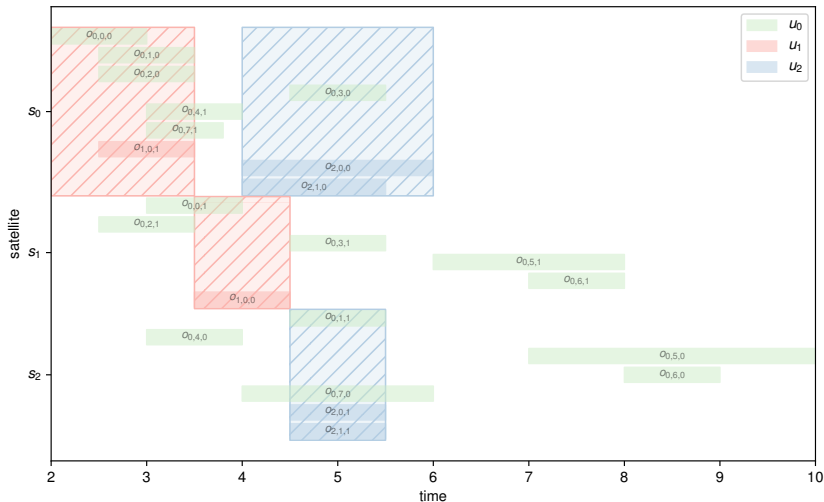
# Inter-Exclusive Coordinated Scheduling

- We focus here on **collective observation scheduling** on a constellation where some users have **exclusive access to some orbit portions**
- ⇒ Answer to strong user expectations to benefit both from a shared system (to reduce costs) and a proprietary system (total control and confidentiality)



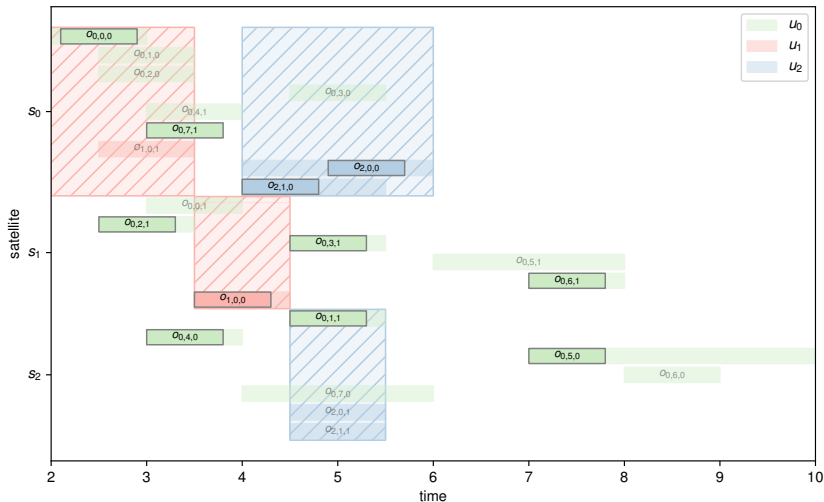
# Scheduling Observations on an EOS Constellation

## Illustrative Example



# Scheduling Observations on an EOS Constellation

## Illustrative Example



# The Problems Behind

- How to **coordinate** exclusive user plans, **without disclosing private plans**, whilst meeting system constraints (memory, energy, etc.)
- How to couple private and non-private observations as to **maximize the system cost-efficiency?**



*Earth Observation Satellite Constellation Scheduling with Exclusives Problem* is a tuple

$$P = \langle \mathcal{S}, \mathcal{U}, \mathcal{R}, \mathcal{O} \rangle$$

- $\mathcal{S} = \{s = \langle t_s^{\text{start}}, t_s^{\text{end}}, \kappa_s, \tau_s \rangle\}$  is a set of satellites
- $\mathcal{U} = \{u = \langle e_u, p_u \rangle\}$  is a set of users
- $\mathcal{R} = \{r = \langle t_r^{\text{start}}, t_r^{\text{end}}, \Delta_r, \rho_r, p_r, u_r, \theta_r \rangle\}$  is a set of requests
- $\mathcal{O} = \{o = \langle t_o^{\text{start}}, t_o^{\text{end}}, \Delta_o, r_o, \rho_o, s_o, u_o, p_o \rangle\}$  is a set of observation opportunities

A *solution* to an EOSCSP is a mapping  $\mathcal{M} = \{(o, t) \mid o \in \mathcal{O}, t \in [t_o^{\text{start}}, t_o^{\text{end}}]\}$

s.t. the overall reward is maximized (sum of the rewards of the scheduled observations):

$$\mathbf{argmax}_{\mathcal{M}} \sum_{(o,t) \in \mathcal{M}} \rho_o$$



# How to Solve EOSCSPs?

---

# How to Solve EOSCSPs?

---

- Centralized allocation

# How to Solve EO SCSPs?

- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up

$$\begin{aligned}
 & \text{maximize}_{x_{i,o}} \sum_{o \in O, i \in I} p_i x_{i,o} & (1) \\
 \text{s.t.} & \\
 & 2 - \beta_{i,o,p} - \beta_{j,o,p} \geq x_{i,o} & \forall i \in S, \forall j, p \in O, o \neq p & (2) \\
 & 2 - \beta_{i,o,p} - \beta_{j,o,p} \geq x_{j,o} & \forall i \in S, \forall j, p \in O, o \neq p & (3) \\
 & \beta_{i,o,p} + \beta_{j,o,p} \leq 3 - x_{i,o} - x_{j,o} & \forall i \in S, \forall j, p \in O, o \neq p & (4) \\
 & \beta_{i,o,p} + \beta_{j,o,p} \leq 1 & \forall i \in S, \forall j, p \in O, o \neq p & (5) \\
 & t_{i,o} - t_{j,o} \geq \tau_i(n_i, p) + \Delta_o - \Delta_{o,p}^{\max} \beta_{j,o,p} & \forall i \in S, \forall j, p \in O, o \neq p, \beta_{j,o,p} > 0 & (6) \\
 & t_{i,o} - t_{j,o} \geq \tau_j(n_j, p) + \Delta_p - \Delta_{p,o}^{\max} \beta_{i,o,p} & \forall i \in S, \forall j, p \in O, o \neq p, \beta_{i,o,p} > 0 & (7) \\
 & \sum_{o \in O} x_{i,o} \leq n_i & \forall i \in S & (8) \\
 & \sum_{i \in I(n_i)} x_{i,o} \leq 1 & \forall o \in R & (9) \\
 & x_{i,o} \in \{0, 1\} & & \\
 & t_{i,o} \in [\frac{c_{i,o}^{\min}}{v_i}, \frac{c_{i,o}^{\max}}{v_i}] \subset \mathbb{R} & \forall i \in S, \forall o \in O & (10) \\
 & \beta_{i,o,p} \in \{0, 1\} & \forall i \in S, \forall p \in O & (11) \\
 & \text{with } \Delta_{i,o,p}^{\max} = \frac{c_{i,o}^{\max}}{v_i} - \frac{c_{j,o}^{\min}}{v_j} + \Delta_o + \tau^i(n_i, p) & \forall i \in S, \forall j, p \in O, o \neq p & (12)
 \end{aligned}$$

# How to Solve EOCSPPs?

- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up
  - Heuristic solving (e.g. greedy)

$$\text{maximize}_{x_{i,p}} \sum_{i \in O, p \in O} A_{i,p} x_{i,p} \quad (1)$$

$$\text{s.t.}$$

$$2 - \beta_{i,p} - \beta_{p,i} \geq x_{i,p} \quad \forall i \in S, \forall p \in O, o \neq p \quad (2)$$

$$2 - \beta_{i,p} - \beta_{p,i} \geq x_{i,p} \quad \forall i \in S, \forall p \in O, o \neq p \quad (3)$$

$$\beta_{i,p} + \beta_{p,i} \leq 3 - x_{i,p} - x_{p,i} \quad \forall i \in S, \forall p \in O, o \neq p \quad (4)$$

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$$t_{i,p} - t_{p,i} \geq \tau_i(n_i, p) + \Delta_{i,p} - \Delta_{i,p}^{\text{max}} \beta_{i,p} \quad \forall i \in S, \forall p \in O, o \neq p \quad (6)$$

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$$\sum_{i \in O} x_{i,i} \leq N_{i,i} \quad \forall i \in S \quad (8)$$

$$x_{i,p} \geq 0 \quad \forall i \in S, \forall p \in O \quad (9)$$

$$x_{i,p} \in \mathbb{R} \quad \forall i \in S, \forall p \in O \quad (10)$$

$$x_{i,p} \in \mathbb{Z} \quad \forall i \in S, \forall p \in O \quad (11)$$

$$x_{i,p} \in \mathbb{N} \quad \forall i \in S, \forall p \in O \quad (12)$$

---

**Algorithm 1: Greedy EOCSPP solver**

**Data:** An EOCSPP  $P = (S, U, R, O)$

**Result:** An assignment  $M$

with

```

M ← {}
Osorted ← sort(O)
R ← {(s, []) | s ∈ S}
for o ∈ Osorted do
    t ← first_slot(o, P, R)
    if t ≠ ∅ then
        M ← M ∪ {(o, t)}
        Osorted ← Osorted \ {o}
return S
    
```

# How to Solve EOCSPPs?

- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up
  - Heuristic solving (e.g. greedy)
  - ✗ private plan disclosure

maximize  $\sum_{i \in O, p \in P} A_i f_{i,p}$  (1)  
 s.t.  
 $2 - \beta_{i,p} - \beta_{j,p} \geq x_{i,p}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (2)  
 $2 - \beta_{i,p} - \beta_{j,p} \geq x_{j,p}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (3)  
 $\beta_{i,p} + \beta_{j,p} \leq 3 - x_{i,p} - x_{j,p}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (4)  
 $\beta_{i,p} + \beta_{j,p} \leq 1$   $\forall i \in S, \forall j, p \in O, o \neq p$  (5)  
 $t_{i,p} - t_{j,p} \geq \tau_i(n_i, p) + \Delta_{i,j} - \Delta_{i,j}^{max} \beta_{i,p}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (6)  
 $t_{i,p} - t_{j,p} \geq \tau_j(n_j, p) + \Delta_{j,i} - \Delta_{j,i}^{max} \beta_{j,p}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (7)  
 $\sum_{i \in O} x_{i,p} \leq N_p$   $\forall p \in S$  (8)  
 $x_{i,p} \geq 0$  (9)  
 $t_{i,p} \in \mathcal{R}$  (10)  
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 $t_{i,p} \in \mathcal{O}$  (12)

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**Result:** An assignment  $\mathcal{M}$

with  $\mathcal{M} \leftarrow \{\}$

$\mathcal{O}_{sorted} \leftarrow \text{sort}(\mathcal{O})$

$R \leftarrow \{(s, \emptyset) \mid s \in S\}$

**for**  $o \in \mathcal{O}_{sorted}$  **do**

$t \leftarrow \text{first\_slot}(o, P, R)$

**if**  $t \neq \emptyset$  **then**

$\mathcal{M} \leftarrow \mathcal{M} \cup \{(o, t)\}$

$\mathcal{O}_{sorted} \leftarrow \mathcal{O}_{sorted} \setminus \{o\}$

**return**  $\mathcal{M}$

# How to Solve EOCSPPs?

- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up
  - Heuristic solving (e.g. greedy)
  - ✗ private plan disclosure
- Distributed allocation

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 s.t.  
 $2 - \beta_{i,o,p} - \beta_{j,o,p} \geq x_{i,o}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (2)  
 $2 - \beta_{i,o,p} - \beta_{j,o,p} \geq x_{j,o}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (3)  
 $\beta_{i,o,p} + \beta_{j,o,p} \leq 3 - x_{i,o} - x_{j,o}$   $\forall i \in S, \forall j, p \in O, o \neq p$  (4)  
 $\beta_{i,o,p} \leq 1$   $\forall i \in S, \forall o, p \in O, o \neq p$  (5)  
 $t_{i,o} - t_{j,o} \geq \tau_i(n_i, p) + \Delta_o - \Delta_{exp}^{max} \beta_{i,o,p}$   $\forall i \in S, \forall j, o, p \in O$  (6)  
 $t_{i,o} - t_{j,o} \geq \tau_j(n_j, o) + \Delta_p - \Delta_{exp}^{max} \beta_{j,o,p}$   $\forall i \in S, \forall j, o, p \in O$  (7)  
 $\sum_{i \in O} x_{i,o} \leq N_o$   $\forall o \in S$  (8)  
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return S
    
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- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up
  - Heuristic solving (e.g. greedy)
    - ✗ private plan disclosure
- Distributed allocation
  - Auctions (e.g. PSI, SSI, CBBA)

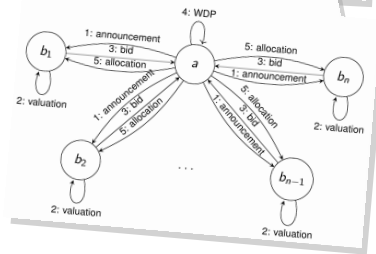
$$\begin{aligned}
 & \text{maximize}_{x_{i,p}} \sum_{i \in O, p \in S} \beta_{i,p} x_{i,p} \\
 & \text{st.} \\
 & 2 - \beta_{i,n,p} - \beta_{i,p,n} \geq x_{i,p} \quad \forall i \in S, \forall n, p \in O, n \neq p \quad (2) \\
 & 2 - \beta_{i,n,p} - \beta_{i,p,n} \geq x_{i,p} \quad \forall i \in S, \forall n, p \in O, n \neq p \quad (3) \\
 & \beta_{i,n,p} + \beta_{i,p,n} \leq 3 - x_{i,p} - \beta_{i,p} \quad \forall i \in S, \forall n, p \in O, n \neq p \quad (4) \\
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 & \sum_{i \in O} x_{i,p} \leq N_p \quad \forall p \in S \quad (8) \\
 & x_{i,p} \geq 0 \quad \forall i \in S, \forall p \in S \quad (9) \\
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 \end{aligned}$$

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**Data:** An EOCSPP  $P = (S, \mathcal{U}, \mathcal{R}, \mathcal{O})$

**Result:** An assignment  $\mathcal{M}$

with  $\mathcal{M} \leftarrow \{ \}$   
 $\text{sorted} \leftarrow \text{sort}(\mathcal{O})$   
 $\forall i \in S$



# How to Solve EOCSPPs?

- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up
  - Heuristic solving (e.g. greedy)
  - ✗ private plan disclosure
- Distributed allocation
  - Auctions (e.g. PSI, SSI, CBBA)
  - Distributed optimization (e.g. DCOPs)

$$\begin{aligned}
 & \text{maximize} && \sum_{i \in O} \alpha_i x_{i,o} \\
 & \text{st.} && \\
 & && 2 - \beta_{i,o,p} - \beta_{j,o,p} \geq \alpha_{i,o} \quad \forall i \in S, \forall j, p \in O, o \neq p \quad (2) \\
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 & && \alpha_{i,o} \geq 0 \quad (10) \\
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## Algorithm 1: Greedy EOCSPP solver

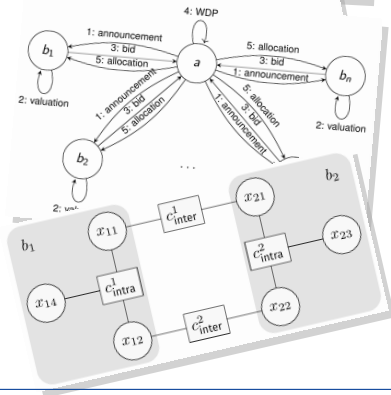
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$\forall i \in S$





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- Centralized allocation
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    - ✓ plans remain private

$$\begin{aligned}
 & \text{maximize} && \sum_{i \in O} \alpha_i x_{i,o} \\
 & \text{st.} && \\
 & && 2 - \beta_{i,o,p} - \beta_{j,o,p} \geq \alpha_{i,o} && \forall i \in S, \forall j, p \in O, o \neq p && (2) \\
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 & && i \in S && && && (10) \\
 & && j \in S && && && (11) \\
 & && o \in S && && && (12)
 \end{aligned}$$

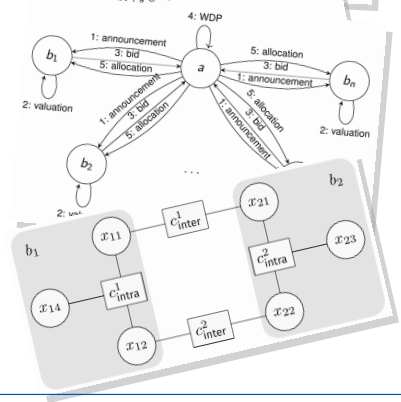
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# How to Solve EOCSPPs?

- Centralized allocation
  - Exact solving (e.g. MILP), but won't scale-up
  - Heuristic solving (e.g. greedy)
    - ✗ private plan disclosure
- Distributed allocation
  - Auctions (e.g. PSI, SSI, CBBA)
  - Distributed optimization (e.g. DCOPs)
    - ✓ plans remain private
    - ⚠ requires some coordination/communication

$$\begin{aligned}
 & \text{maximize} && \sum_{i \in O} \alpha_i x_{i,p} \\
 & \text{st.} && \\
 & && 2 - \beta_{i,p} - \beta_{j,p} \geq \alpha_{i,p} && \forall i \in S, \forall j, p \in O, i \neq j && (2) \\
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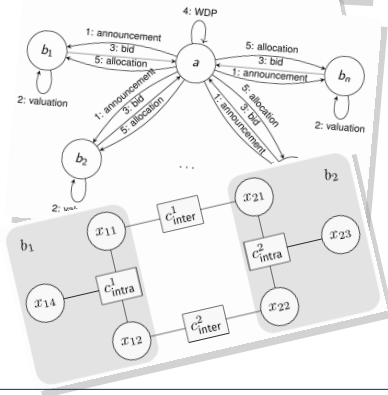
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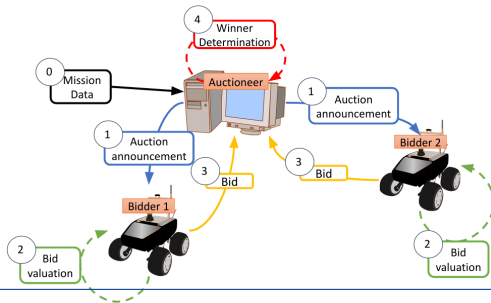


# Auction-based Coordination for EOSCSP

## Focus on Resource/Task Allocation

Many application fields, as Collective Robotics, make use of market-based approach to allocate tasks/resources to robots

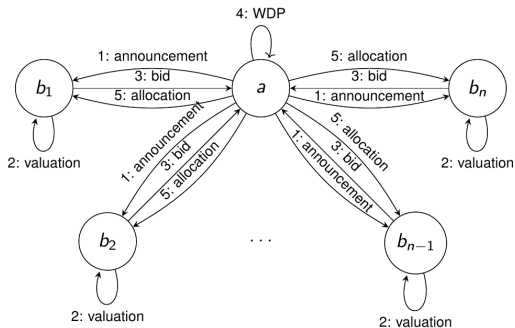
- A set of **resources** (robots, satellites, etc.),  $R = \{r_1, \dots, r_{|R|}\}$
  - A set of **tasks**,  $T = \{t_1, \dots, t_{|T|}\}$ , each having a time-related and operation constraints
  - Find an allocation of tasks to resources, wrt. some consistency constraints
- ≈ **multi-item allocation**: each resource is allocated several tasks (bundle)



# Auction-based Coordination for EOSCSP

Allocating non exclusive observations to best exclusive portions

Auction-based approaches are relevant for satellite task allocation [PHILLIPS and PARRA, 2021]

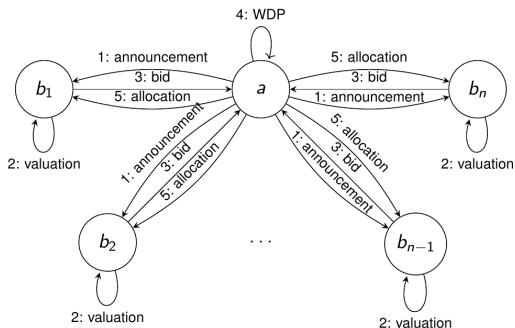


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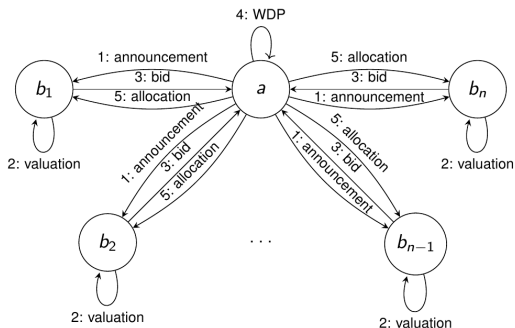


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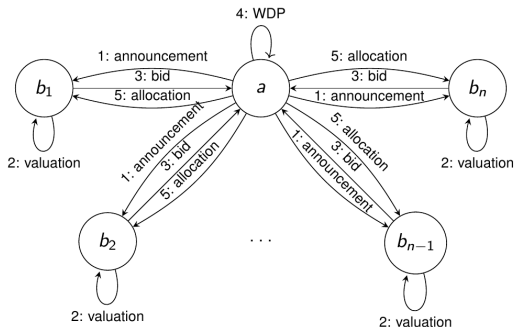
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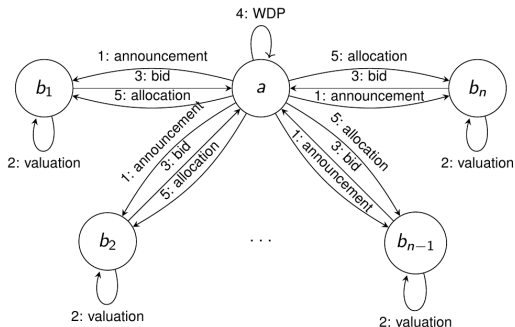
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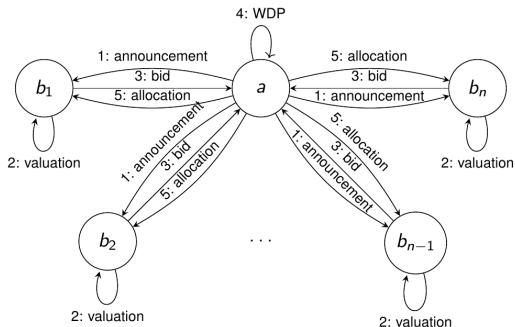
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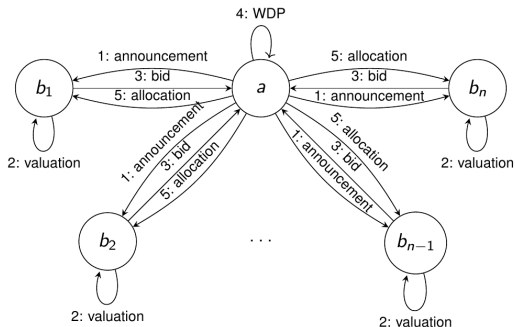


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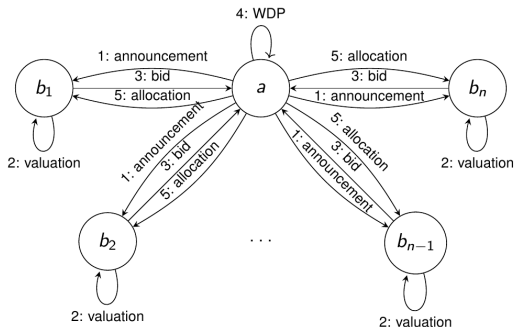


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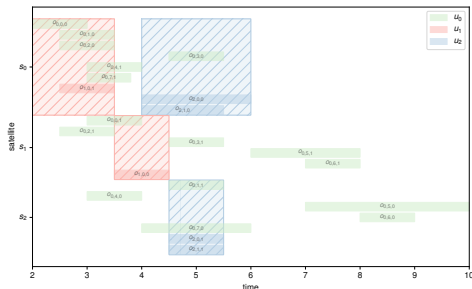


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- **Consensus-based Bundle Auction (CBBA)** [CHOI et al., 2009]
  - Each agent bids on some *bundle of tasks* and *converge to a consensus* with other agents

# Applying Auction-based Allocation to EOSCSP

## General Scheme

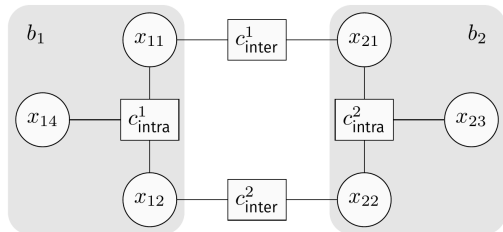
- 1 Identify non exclusive requests possibly fulfilled in exclusive portions
- 2 Send identified requests to exclusive users
- 3 Solve the allocation problem using PSI, SSI or CBBA
  - Bids are computed as the **best marginal costs** of integrating requests in their current plans (which amounts to solve scheduling problems...)
- 4 Allocate as many remaining requests outside exclusive windows



# DCOP-based Coordination for EOSCSP

Allocating non exclusive observations to best exclusive portions

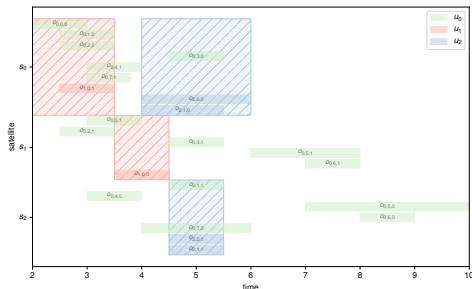
- Consider the **collective decision** for allocating non exclusive tasks to exclusive windows
- Collective decision to coordinate exclusive users' decisions modeled as a **distributed constraint optimization problem** (DCOP)
- As for auctions, exclusive users aim to **minimizing the marginal cost** of integrating non exclusive tasks in their schedule, while meeting some operational constraints



# DCOP-based Coordination for EOSCSP

## General Scheme

- 1 Identify non exclusive requests possibly fulfilled in exclusive windows
- 2 Send each identified request  $r$  to exclusives users, one by one
- 3 Solve the problem of  $r$  using a DCOP solution method (e.g. DPOP [Petcu2005])
  - Costs are computed as the **best marginal cost** of integrating requests in their current plan (which amounts to solve a scheduling problem...)
- 4 Allocate as many remaining requests outside exclusive windows



# DCOP-based Coordination for EOSCSP

## DCOP Model

---

A DCOP  $\langle \mathcal{A}, \mathcal{X}, \mathcal{D}, \mathcal{C}, \mu \rangle$  is defined for a given request  $r$ , and a current scheduling

# DCOP-based Coordination for EOSCSP

## DCOP Model

A DCOP  $\langle \mathcal{A}, \mathcal{X}, \mathcal{D}, \mathcal{C}, \mu \rangle$  is defined for a given request  $r$ , and a current scheduling

- The agents are the exclusive users which can potentially schedule  $r$ :

$$\mathcal{A} = \{u \in \mathcal{U}^{\text{ex}} \mid \exists (s, (t_u^{\text{start}}, t_u^{\text{end}})) \in e_u, \exists o \in \theta_r \text{ s.t. } s_o = s, [t_u^{\text{start}}, t_u^{\text{end}}] \cap [t_o^{\text{start}}, t_o^{\text{end}}] \neq \emptyset\} \quad (1)$$



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- Each agent  $u$  owns binary decision variables, one for each observation  $o \in \mathcal{O}[u]^r$  and exclusive  $e$  in its exclusives  $e_u$ , stating whether it schedules  $o$  in  $e$  or not:

$$\mathcal{X} = \{x_{e,o} \mid e \in \bigcup_{u \in \mathcal{A}} e_u, o \in \mathcal{O}[u]^r\} \quad (2)$$

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- $\mu$  associates each variable  $x_{e,o}$  to  $e$ 's owner

# DCOP-based Coordination for EOSCSP (cont.)

## DCOP Model

- Constraints should check that at most one observation is scheduled per request (4), that satellites are not overloaded (5), that at most one agent serves the same observation (6)

$$\sum_{e \in \bigcup_{u \in \mathcal{A}} e_u} x_{e,o} \leq 1, \quad \forall u \in \mathcal{X}, \forall o \in \mathcal{O}[u]^r \quad (4)$$

$$\sum_{o \in \{o \in \mathcal{O}[u]^r \mid u \in \mathcal{A}, s_o = s\}, e \in \bigcup_{u \in \mathcal{A}} e_u} x_{e,o} \leq \kappa_s^*, \quad \forall s \in \mathcal{S} \quad (5)$$

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- The cost to integrate an observation in the current user's schedule should be assessed to guide the optimization process

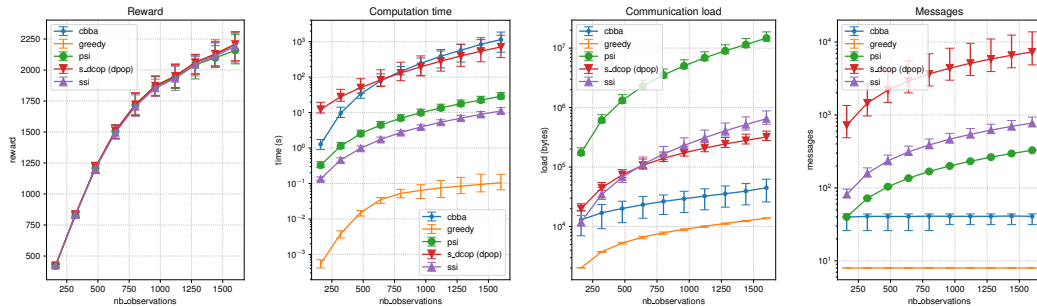
$$c(x_{e,o}) = \pi(o, \mathcal{M}_{u_o}), \quad \forall x_{e,o} \in \mathcal{X} \quad (7)$$

where  $\pi$  evaluates the best cost obtained when scheduling  $o$  and any combination of observations from  $\mathcal{M}_{u_o}$ , as to consider all possible revisions of  $u_o$ 's current schedule

$$\mathcal{C} = \{(4), (5), (6), (7)\} \quad (8)$$

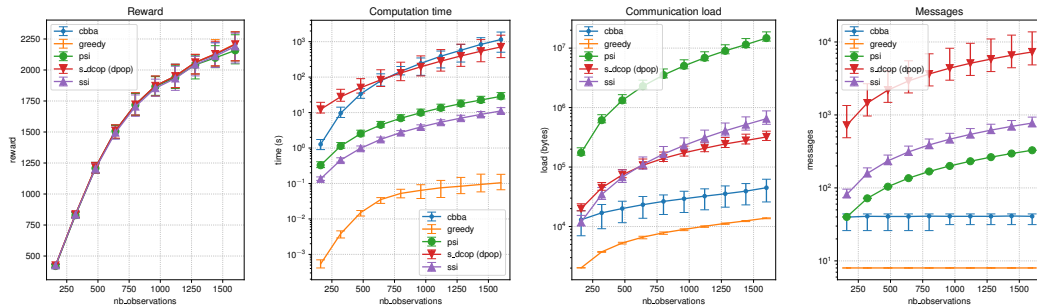
# Highly conflicting randomly generated problems

5-min horizon with overlapping requests and limited capacity



# Highly conflicting randomly generated problems

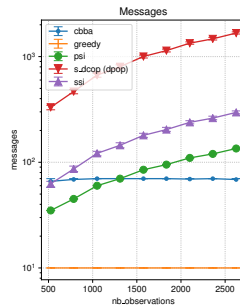
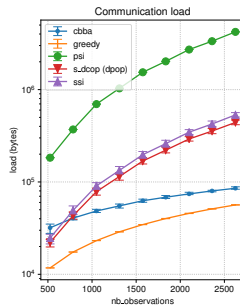
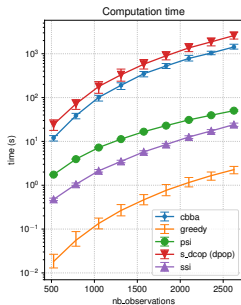
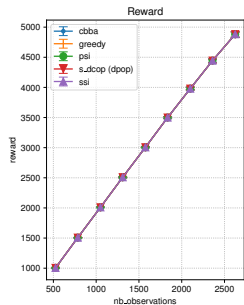
5-min horizon with overlapping requests and limited capacity



- ✗ cbba and s\_dcop requires extra-computation time ( $\approx 1000s$ )
- ✓ cbba and s\_dcop provide the best solutions wrt. reward
- ✓ cbba exchanges fewer messages of small size
- ✓ ssi remains the best compromise wrt. solution quality, computation time and communication load

# Realistic randomly generated problems

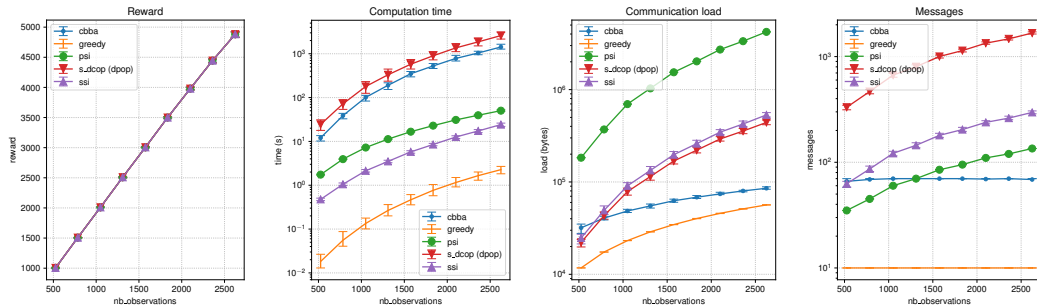
6-hour horizon with numerous requests and large capacity





# Realistic randomly generated problems

6-hour horizon with numerous requests and large capacity



- ✓ cbba does require less time to compute than s\_dcop
- ✓ s\_dcop and cbba can perform many computation concurrently
- ⇒ There is room for computation speedup in real distributed settings

# Where to find detailed info?

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- Initial model definition [PICARD, 2022]
- Auction-based and DCOP-based solution methods [ibid.]
- More complex requests and decentralized auctions [PICARD, 2023a]
- Some data [PICARD, 2023b]

# Outline

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- 1 Introduction
- 2 Challenges in Earth Observation Constellation Operations
- 3 Focus #1: Sharing Space Assets
- 4 Focus #2: Coordinating Asset Usage
- 5 Conclusion

# Wrap-up

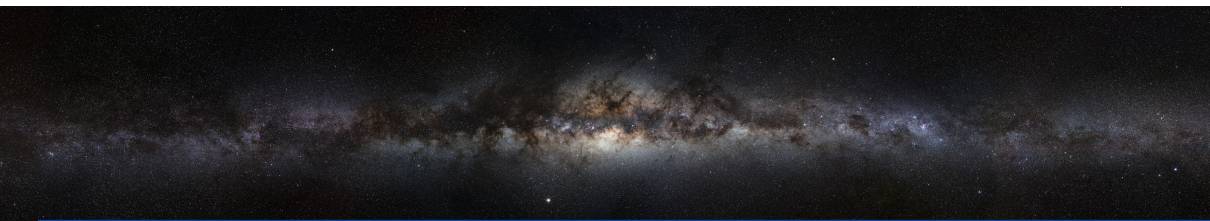
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- Key terms for NewSpace: multi-asset, multi-user, multi-system...
- Asset sharing means **cost-efficiency**, but requires **automated coordination** and **privacy/sovereignty** preservation

# Wrap-up

---

- How to coordinate such composite systems?
  - Efficiency
  - Fairness
  - Explainability



# Wrap-up

---

- How to coordinate such composite systems?
  - Efficiency
  - Fairness
  - Explainability
- Multi-agent Systems
  - Resource allocation and combinatorial auctions
  - Distributed optimization
  - Federated and multi-agent learning
  - ...

# Our Next Steps

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- Even more complex requests
  - Periodic intra-/inter-day, short-/long-term
  - Large area and responsiveness
- Even more complex systems
  - Weather uncertainties
  - Constellation federations



# Acknowledgements

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Jean-Loup Farges, Cédric Pralet, Stéphanie Roussel, and Sara Maqrot  
(at ONERA)

Clément Caron and Jonathan Guerra  
(at ADS)



# Acknowledgements

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This work has been performed with the support of the French government in the context of the “Programme d’Investissements d’Avenir”, namely by the BPI PSCC project “LiChIE”






**bpi**france

**Thank you for your attention!**  
**Any question?**

[www.onera.fr](http://www.onera.fr)

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